The Panpositionable Hamiltonicity of k-ary n-cubes

Chia-Jung Tsai and Shin-Shin Kao

Abstract—The hypercube Q_n is one of the most well-known and popular interconnection networks and the k-ary n-cube Q_n^k is an enlarged family from Q_n that keeps many pleasing properties from hypercubes. In this article, we study the panpositionable hamiltonicity of Q_n^k for $k \geq 3$ and $n \geq 2$. Let x, y of $V(Q_n^k)$ be two arbitrary vertices and $\mathcal C$ be a hamiltonian cycle of Q_n^k . We use $d_{\mathcal C}(x,y)$ to denote the distance between x and y on the hamiltonian cycle $\mathcal C$. Define l as an integer satisfying $d(x,y) \leq l \leq \frac{1}{2}|V(Q_n^k)|$. We prove the followings:

- When k=3 and $n\geq 2$, there exists a hamiltonian cycle C of Q_n^k such that $d_C(x,y)=l$.
- When $k \geq 5$ is odd and $n \geq 2$, we request that $l \notin S$ where S is a set of specific integers. Then there exists a hamiltonian cycle C of Q_n^k such that $d_C(x,y) = l$.
- When $k \geq 4$ is even and $n \geq 2$, we request l-d(x,y) to be even. Then there exists a hamiltonian cycle C of Q_n^k such that $d_C(x,y)=l$.

The result is optimal since the restrictions on l is due to the structure of \mathcal{Q}_n^k by definition.

Index Terms—Hamiltonian, panpositionable, bipanpositionable, k-ary n-cube.

I. Introduction

THE n-dimensional hypercube Q_n is one of the most well-known and popular interconnection networks due to its excellent properties as the following: it is vertex-symmetric and edge-symmetric; it is hamiltonian; it allows cycle/path embedding when faults occur and so on. (See [1], [2] for these results and their references). Therefore, numerous studies have been devoted to the hypercube family [3]–[6], [11], [12].

The k-ary n-cube Q_n^k is an enlarged family from Q_n that keeps many pleasing properties from hypercubes. More precisely, each vertex of Q_n^k is labeled by a n-bit finite sequence $(u_0,u_1,...,u_{n-1})$, where $0 \le u_i \le k-1$ for $0 \le i \le n-1$, and every two vertices u and v are adjacent if and only if $|u_i-v_i|=1$ or k-1 for some i and $u_j=v_j$ for any $0 \le j \le n-1$ with $j \ne i$. It is obviously that the hypercube Q_n is indeed a subclass of the k-ary n-cube when k=2. Some properties of Q_n^k mentioned in [6] are listed here: it is known that Q_n^k is vertex-symmetric and edge-symmetric [3]; it is hamiltonian [4], [5]; it has diameter $n\lfloor \frac{k}{2} \rfloor$ [4], [5]; it has a recursive structure; and it contains many important interconnection networks such as cycles (of certain lengths) [3], meshes (of certain dimensions) [4], and even hypercubes

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(of certain dimensions) [5]. However, as opposed to Q_n , Q_n^k has not received enough attention. In this article, we want to prove the panpositionability of Q_n^k . Readers can refer to [7] for the concept of panpositionability. A hamiltonian graph G is panpositionable if for any two different vertices u and v of G and any integer l with $d_G(u,v) \leq l \leq \frac{|V(G)|}{2}$, there exists a hamiltonian cycle C of G with $d_C(u,v) = l$. Similar to the hamiltonicity for the communication between processors in an interconnection network, panpositionable hamiltonicity allows more flexible communication in a hamiltonian network. It is easy to see that the panpositionable hamiltonian property inherits the hamiltonian property and advances it further [8].

The article is organized as follows. In Section 2, we introduce the graph terminologies and notations used in this paper, the precise definition of Q_n^k , and two lemmas. In Section 3, we study the panpositionability of Q_n^k , where $k \geq 3$ is an odd integer and $n \geq 2$ is an integer. In Section 4, we study the panpositionability in the bipartite version of Q_n^k , where $k \geq 4$ is an even integer and $n \geq 2$ is an integer. Our conclusion is given in the last section.

II. PRELIMINARIES

For the graph definitions and notations we follow [9]. G = (V, E) is a graph if V is a finite set and E is a subset of $\{\{u,v\}|\{u,v\}$ is an unordered pair of $V\}$. We say that V is the vertex set and E is the edge set of G. Two vertices u and v are adjacent if $\{u,v\} \in E$. A path is represented by a finite sequence of vertices $\langle v_0, v_1, v_2, ..., v_n \rangle$, where every two consecutive vertices are adjacent. If P is a path represented by $\langle v_0, v_1, v_2, ..., v_n \rangle$, then we define $\operatorname{inv}(P) = \langle v_n, v_{n-1}, v_{n-2}, ..., v_0 \rangle$. The *length* of a path P is the number of edges in P. We write the path $\langle v_0, v_1, ..., v_n \rangle$ as $\langle v_0, v_1, ..., v_{s-1}, P_1, v_{i+1}, ..., v_{j-1}, P_2, v_{t+1}, ..., v_n \rangle$, where $P_1 = \langle v_s, v_{s+1}, ..., v_i \rangle$ and $P_2 = \langle v_j, v_{j+1}, ..., v_t \rangle$. We use $d_G(u,v)$ to denote the distance between u and v in G, i.e., the length of the shortest path between u and v in G. A cycle is a path of at least three vertices such that the first vertex is the same as the last vertex. A hamiltonian cycle of G is a cycle that visits every vertex of G exactly once. We use $d_C(u,v)$ to denote the distance between u and v in a cycle C of G, i.e., the length of the shorter path between u and v in C. A hamiltonian graph is a graph with a hamiltonian cycle.

A hamiltonian path in a graph G is a path joining two distinct vertices u and v of G that visits every vertex of G exactly once. A graph G is hamiltonian-connected if there is a hamiltonian path joining any two distinct vertices of G. Note that any (nontrivial) bipartite graph cannot be hamiltonian-connected, whereas a bipartite graph is hamiltonian laceable if there exists a hamiltonian path joining every two vertices which are in distinct partite [10].

The concept of hamiltonian panpositionability was first proposed by S. Kao etc. [7]. A hamiltonian graph G is panpositionable if for any two different vertices u and v of G and any integer l with $d_G(u,v) \leq l \leq \frac{|V(G)|}{2}$, there exists a hamiltonian cycle C of G with $d_C(u,v) = l$. A graph $G = (V_0 \cup V_1, E)$ is bipartite if $V(G) = V_0 \cup V_1$ and E(G) is a subset of $\{\{u,v\}|u \in V_0, v \in V_1\}$. A hamiltonian bipartite graph G is bipanpositionable if for any two different vertices u and v of G and any integer l with $d_G(u,v) \leq l \leq \frac{|V(G)|}{2}$ and $(l-d_G(u,v))$ is even, there exists a hamiltonian cycle C of G with $d_C(u,v) = l$.

The k-ary n-cube, Q_n^k is defined for all integers $k \geq 2$ and $n \geq 1$. The subclass Q_n^2 is the well-studied hypercube family. The subclass Q_1^k with $k \geq 3$ is defined as the cycle of length k. The k-ary n-cube, Q_n^k , for $k \geq 3$ and $n \geq 2$ is defined as follows. Let $u \in V(Q_n^k)$ be represented by $(u_0, u_1, ..., u_{n-1})$, where $0 \leq u_i \leq k-1$. u and v are adjacent if and only if $|u_i-v_i|=1$ or k-1 for some i and $u_j=v_j$ for any $0 \leq j \leq n-1$ with $j \neq i$. It is shown that Q_n^k is bipartite if k is even [11]. Here we mention some properties of Q_n^k that will be used in this article.

It is known that Q_n^k is vertex-symmetric and edge-symmetric. Moreover, given any two distinct vertices (u_1,u_2) and (v_1,v_2) of Q_2^k , there is an automorphism of Q_2^k mapping (u_1,u_2) and (v_1,v_2) to (m,0) and (0,n). Each vertex of Q_n^k is represented by a n-bit tuple, and we will call the dth-bit the dth dimension. We can partition Q_n^k over dimension d by fixing the dth element of any vertex tuple at some value a, for every $a \in \{0,1,\dots,k-1\}$. This results in k copies $Q_{d,n-1}^{k,0}, Q_{d,n-1}^{k,1}, \dots, Q_{d,n-1}^{k,k-1}$ of Q_{n-1}^k , with corresponding vertices in $Q_{d,n-1}^{k,0}, Q_{d,n-1}^{k,1}, \dots, Q_{d,n-1}^{k,k-1}$ joined in a cycle of length k (in dimension d) [6]. It is proven in [11], [12] that Q_n^k is hamiltonian connected for odd k and Q_n^k is hamiltonian laceable for even k.

Note that the length of a path between u and v in Q_n^k , where $k \geq 5$ is an odd integer, can not be arbitrary. For example, in Q_2^5 , for any two vertices u and v and d(u,v)=1, there exists no path P between u and v with |P| = 2. In fact, we have the following observation. Given two vertices u = $(u_0, u_1, ..., u_{n-1})$ and $v = (v_0, v_1, ..., v_{n-1})$ of Q_n^k . Define the number $m_i = \min\{|u_i - v_i|, k - |u_i - v_i|\}$, where $0 \le i \le n - 1$. Let $s = \max\{m_i : 0 \le i \le n-1\}$. Then there exists no path between u and v with length r = d(u, v) - s + k - s - 2l =d(u,v)+k-2s-2l, where l is an integer and $1 \le l \le \frac{k}{2}-s$. Consequently, we modify the concept of panpositionability of Q_n^k by saying that Q_n^k is nearly-panpositionable if for any two distinct vertices x and y of Q_n^k and for any integer l' with $d(x,y) \leq l' \leq \frac{|V(Q_n^k)|}{2}$ and $l' \notin \{r: r = d(u,v) + k - 2s - 2l \text{ for } 1 \leq l \leq \frac{k}{2} - s\}$, there exists a hamiltonian cycle C of Q_n^k with $d_C(x,y) = l'$. Therefore, in this article, we want to prove that Q_n^3 is panpositionable, Q_n^k is nearly-panpositionable if $k \geq 5$ is an odd integer, and is bipanpositionable if $k \geq 4$ is an even integer. First of all, we prove the following two

Lemma 1. Let k be an integer with $k \geq 3$. For any path P with length 2 in Q_2^k , there exists a hamiltonian cycle of Q_2^k that contains P.

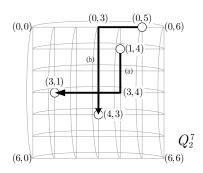


Fig. 1. (a) $f_{-3}^2(1,4)$ and (b) $h_{-2}^4(0,5)$.

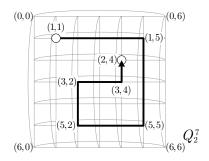


Fig. 2. $H^{\vec{a}}_{\vec{b},\vec{3}}(1,1)$, where $\vec{a}=(4,-2,-1)$ and $\vec{b}=(4,-3,2)$.

Proof: Let c,r,i be nonzero integers, $\frac{c}{|c|}=s$, $\frac{r}{|r|}=t$, $\vec{a}=(a_1,a_2,...,a_i)$ and $\vec{b}=(b_1,b_2,...,b_i)$. If c=0, then s=0. Similarly, if r=0, then t=0. To construct the required hamiltonian cycles, we define some path patterns in the following.

 $\begin{array}{ll} f_r^c(x,y) &= \langle (x,y), (x+s\cdot 1,y), (x+s\cdot 2,y), ..., (x+c,y), (x+c,y+t\cdot 1), (x+c,y+t\cdot 2), ..., (x+c,y+r) \rangle; \\ h_b^c(x,y) &= \langle f_p^0(x,y), f_0^c(x,y+r) \rangle; \\ H_{\overline{b},i}^{\overline{a}}(x,y) &= \langle h_{b_1}^{a_1}(x,y), h_{b_2}^{a_2}(x+a_1,y+b_1), h_{b_3}^{a_3}(x+a_1+a_2,y+b_1+b_2), ..., h_{b_i}^{a_i}(x+\sum_{n=1}^{i-1}a_n,y+\sum_{n=1}^{i-1}b_n) \rangle. \end{array}$

Please see Fig. 1 and Fig. 2 for an illustration. Fig. 1 is examples of $f_{-3}^2(1,4)$ and $h_{-2}^4(0,5)$. Note that $f_{-3}^2(1,4)=\langle (1,4),(2,4),(3,4),(3,3),(3,2),(3,1)\rangle$ and $h_{-2}^4(0,5)=\langle f_{-2}^0(0,5),f_0^4(0,3)\rangle=\langle (0,5),(0,4),(0,3),(1,3),(2,3),(3,3),(4,3)\rangle$. Fig. 2 is an example of $H^{\vec{a}}_{\vec{b},3}(1,1)$, where $\vec{a}=(4,-2,-1)$ and $\vec{b}=(4,-3,2)$. Note that $H^{\vec{a}}_{\vec{b},3}(1,1)=\langle h_4^4(1,1),h_{-3}^{-2}(5,5),h_2^{-1}(3,2)\rangle=\langle f_4^0(1,1),f_0^4(1,5),f_{-3}^0(5,5),f_0^{-2}(5,2),f_2^0(3,2),f_0^{-1}(3,4)\rangle=\langle (1,1),(1,2),(1,3),(1,4),(1,5),(2,5),(3,5),(4,5),(5,5),(5,4),(5,3),(5,2),(4,2),(3,2),(3,3),(3,4),(2,4)\rangle$. Let $P=\langle u,x,v\rangle$, where $u=(u_1,u_2)$ and $v=(v_1,v_2)$ in Q_2^k . We have following cases.

Case 1. k is odd.

Case 1.1. Either $u_1 = v_1$ or $u_2 = v_2$. W.L.O.G., let $u = (0,0), \ v = (2,0)$ and $P = \langle u, (1,0), v \rangle$. Let $a_i = (-1)^i (2-k)$, for $i \leq k-1$ and $a_k = 0$; $\vec{b} = (0,-1,-1,...,-1)$. There exists a hamiltonian cycle $C = \langle (0,0), P, (2,0), f_{k-1}^{k-3}(2,0), H_{\vec{b},k}^{\vec{a}}(0,k-1) \rangle$

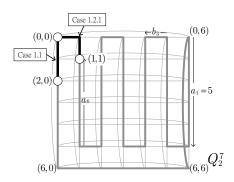


Fig. 3. Examples of Case 1.1 and Case 1.2.1 for k = 7.

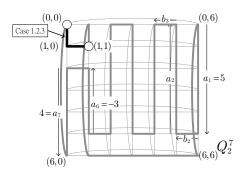


Fig. 4. An Example of Case 1.2.2 for k = 7.

1), $(0,0)\rangle$. Please see Fig. 3 for an illustration. The hamiltonian cycle in Fig. 3 is $C=(0,0),P,(2,0),f_6^4(2,0),H_{\tilde{b},7}^{\vec{a}}(0,6),(0,0)\rangle$ and $H_{\tilde{b},7}^{\vec{a}}(0,6)=(h_0^5(0,6),h_{-1}^{-5}(5,6),h_{-1}^{5}(0,5),h_{-1}^{-5}(5,4),h_{-1}^{5}(0,3),h_{-1}^{-5}(5,2),h_{-1}^{0}(0,1)\rangle$.

Case 1.2. $u_1 \neq v_1, u_2 \neq v_2$. W.L.O.G., let u = (0,0) and v = (1,1).

Case 1.2.1. $P = \langle u, (0, 1), v \rangle$, where $k \ge 3$.

Case 1.2.3. $P = \langle u, (1,0), v \rangle$, where $k \geq 5$.

The hamiltonian cycle is the same as in Case 1.1. Please see Fig. 3 for an illustration.

 $\begin{array}{lll} \overline{\text{Let }a_i = (-1)^i(2-k)}, \text{ for } i \leq k-2, \ a_{k-1} = 4-k \\ \text{and } a_k = k-3; \ \vec{b} = (0,-1,-1,...,-1). \ \text{There exists} \\ \text{a hamiltonian cycle } C = \langle (0,0),P,(0,1),f_{k-2}^0(k-1,1),H_{\vec{b},k}^{\vec{a}}(0,k-1),(0,0)\rangle. \ \text{Please see Fig. 4 for an illustration. The hamiltonian cycle in Fig. 4 is } C = \langle (0,0),P,(0,1),f_5^0(6,1),H_{\vec{b},7}^{\vec{a}}(0,6),(0,0)\rangle \ \text{and } H_{\vec{b},7}^{\vec{a}}(0,6) = \langle h_0^5(0,6),h_{-1}^{-5}(5,6),h_{-1}^5(0,5),h_{-1}^{-5}(5,4),h_{-1}^5(0,3),h_{-1}^{-3}(5,2),h_{-1}^4(2,1)\rangle. \end{array}$

Case 2. k is even.

Case 2.1. Either $u_1 = v_1$ or $u_2 = v_2$. W.L.O.G., let u = (0,0) and v = (2,0) and $P = \langle u, (1,0), v \rangle$.

Let $a_i = (-1)^i(2-k)$, for $3 \le i \le k$, $a_1 = k-3$, $a_2 = 1-k$ and $a_{k+1} = 0$;

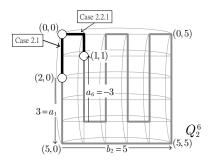


Fig. 5. Examples of Case 2.1 and Case 2.2.1 for k = 6.

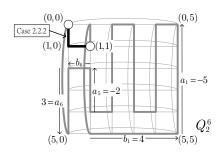


Fig. 6. An Example of Case 2.2.2 for k = 6.

 $\begin{array}{ll} \vec{b} &= (0,k-1,-1,-1,...,-1). \text{ There exists a hamiltonian cycle } C &= \langle (0,0),P,(2,0),H^{\vec{a}}_{\vec{b},k+1}(2,0),(0,0)\rangle. \text{ Please see } \\ \text{Fig. 5 for an illustration. The hamiltonian cycle in Fig. 5} \\ \text{is } C &= \langle (0,0),P,(2,0),H^{\vec{a}}_{\vec{b},7}(2,0),(0,0)\rangle \text{ and } H^{\vec{a}}_{\vec{b},7}(2,0) &= \langle h^3_0(2,0),h^{-5}_5(5,0),h^4_{-1}(0,5),h^{-4}_{-1}(4,4),h^4_{-1}(0,3),h^{-4}_{-1}(4,2),h^0_{-1}(0,1)\rangle. \end{array}$

Case 2.2. $u_1 \neq v_1, u_2 \neq v_2$. W.L.O.G., let u = (0,0) and v = (1,1).

Case 2.2.1. $P = \langle u, (0, 1), v \rangle$

The hamiltonian cycle is the same as in Case 2.1. Please see Fig. 5 for an illustration.

Case 2.2.2. $P = \langle u, (1,0), v \rangle$

The lemma is proved.

To facilitate our derivation in the following, we define some path patterns. We shall use $x_0^i, x_1^i, x_2^i, ..., x_{k^{n-1}-1}^i$ to denote the k^{n-1} vertices of $Q_{d,n-1}^{k,i}$ for some d. For simplicity, denote $Q_{d,n-1}^{k,i}$ as $Q_{n-1}^{k,i}$. Let the path $p(x_a^i, x_b^i) = \langle x_a^i, x_{a_1}^i, x_{a_2}^i, ..., x_b^i \rangle$ and $a_i = (a+i \bmod k^{n-1})$. For example, if $k^{n-1} = 64$, then $p(x_{60}^i, x_2^i) = \langle x_{60}^i, x_{61}^i, x_{62}^i, x_{63}^i, x_1^i, x_2^i \rangle$. It is known that there exists a hamiltonian cycle in Q_{n-1}^k [4]. Thus x_a^i and x_{a+1}^i are adjacent and so are x_a^i and x_a^{i+1} .

Lemma 2. Let k be an integer with $k \geq 3$. For any path P

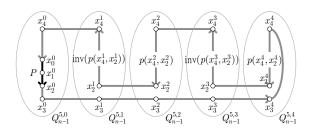


Fig. 7. An Example of Case 1 with k = 5.

with length 2 in Q_n^k , there exists a hamiltonian cycle of Q_n^k that contains P.

Proof: The lemma will be proved by mathematical induction. By Lemma 1, the statement holds for Q_2^k . Using the induction hypothesis, we assume that the statement holds for Q_{n-1}^k , where $n \geq 3$. Now we want to prove that the lemma is true for Q_n^k . There are three cases.

Case 1. P is in $Q_{n-1}^{k,i}$. W.L.O.G., let i=0.

By the induction hypothesis, there exists a hamiltonian cycle By the induction hypothesis, there exists a manifestance C^0 of $Q_{n-1}^{k,0}$ that contains P. Let $P=\langle x_0^0,x_1^0,x_2^0\rangle$ and $C^0=\langle x_0^0,P,x_2^0,x_3^0,...,x_{k^{n-1}-1}^0,x_0^0\rangle$. Since Q_{n-1}^k is hamiltonian [4], let the hamiltonian cycles in Q_{n-1}^k be $C^i=0$ $\langle x_0^i, x_1^i, x_2^i, x_3^i, ..., x_{k^{n-1}-1}^i, x_0^i \rangle.$

- 1) k is odd. Then the hamiltonian cycle is
 $$\begin{split} C &= \langle x_0^0, P, x_0^0, x_3^0, x_3^1, x_3^2, ..., x_4^{k-1}, p(x_4^{k-1}, x_2^{k-1}), \\ &\operatorname{inv}(p(x_4^{k-2}, x_2^{k-2})), p(x_4^{k-3}, x_2^{k-3}), \operatorname{inv}(p(x_4^{k-4}, x_2^k, x_2^k)), \\ &\ldots, p(x_4^2, x_2^2), \operatorname{inv}(p(x_4^1, x_2^1)), p(x_4^0, x_0^0), x_0^0 \rangle. \end{split}$$
- 2) k is even. Then the hamiltonian cycle is $\begin{array}{l} C = \langle x_0^0, P, x_2^0, x_3^0, x_3^1, x_3^2, ..., x_3^{k-1}, \mathrm{inv}(p(x_4^{k-1}, x_2^{k-1})), \\ p(x_4^{k-2}, x_2^{k-2}), \mathrm{inv}(p(x_4^{k-3}, x_2^{k-3})), p(x_4^{k-4}, x_2^{k-4}), ..., \\ p(x_4^2, x_2^2), \mathrm{inv}(p(x_4^1, x_2^1)), p(x_4^0, x_0^0), x_0^0 \rangle. \end{array}$

Please see Fig. 7 for an illustration, where the hamiltonian cycle in Fig. 7 is $C=\langle x_0^0, P, x_2^0, x_3^0, x_3^1, x_3^2, x_3^3, x_3^4, p(x_4^4, x_2^4), \\ \operatorname{inv}(p(x_4^3, x_2^3)), p(x_4^2, x_2^2), \\ \operatorname{inv}(p(x_4^1, x_2^1)), p(x_4^0, x_0^0), x_0^0 \rangle.$

Case 2. P passes through two $Q_{n-1}^{k,i}$. W.L.O.G., those two are $Q_{n-1}^{k,0}$ and $Q_{n-1}^{k,1}$. Let $P=\langle x_0^0,x_1^0,x_1^1\rangle$. In [11], [12], it has been shown that

there exists a hamiltonian path $\langle x_1^i, p(x_1^i, x_0^i), x_0^i \rangle$ in $Q_{n-1}^{k,i}$.

- 1) k is odd. Then the hamiltonian cycle is $\begin{array}{l} C = \langle x_0^0, P, x_1^1, x_1^2, x_1^3, x_1^4, ..., x_1^{k-1}, p(x_2^{k-1}, x_0^{k-1}), \\ \operatorname{inv}(p(x_2^{k-2}, x_0^{k-2})), p(x_2^{k-3}, x_0^{k-3}), \operatorname{inv}(p(x_2^{k-4}, x_0^{k}), x_0^{k-1}), \end{array}$..., $p(x_2^2, x_0^2)$, inv $(p(x_2^1, x_0^1))$, $p(x_2^0, x_0^0)$, x_0^0
- 2) k is even. Then the hamiltonian cycle is $\begin{array}{l} C = \langle x_0^0, P, x_1^1, x_1^2, x_1^3, x_1^4, ..., x_1^{k-1}, \operatorname{inv}(p(x_2^{k-1}, x_0^{k-1})), \\ p(x_2^{k-2}, x_0^{k-2}), \operatorname{inv}(p(x_2^{k-3}, x_0^{k-3})), p(x_2^{k-4}, x_0^{k-4}), ..., \\ p(x_2^2, x_0^2), \operatorname{inv}(p(x_2^1, x_0^1)), p(x_2^0, x_0^0), x_0^0 \rangle. \end{array}$

Please see Fig. 8 for an illustration, where the hamiltonian cycle in Fig. 8 is $C=\langle x_0^0,P,x_1^1,x_1^2,x_1^3, \mathrm{inv}(p(x_2^3,x_0^3)), p(x_2^2,x_0^2), \mathrm{inv}(p(x_2^1,x_0^1)), p(x_2^0,x_0^0), x_0^0 \rangle.$

Case 3. P passes through three $Q_{n-1}^{k,i}$. It is known that we can partition Q_n^k over dimension d by fixing the dth element of any vertex tuple at some value a, for

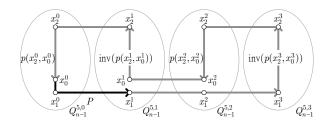


Fig. 8. An Example of Case 2 with k = 4.

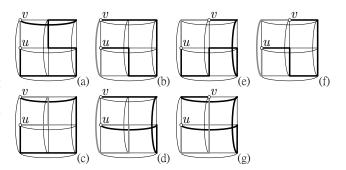


Fig. 9. Illustrations of Lemma 3.

every $a\in\{0,1,...,k-1\}$. In this case, $P=\langle u,x,v\rangle$ passes through three $Q_{n-1}^{k,i}$, i.e., u,x and v have the same value in at least one element of vertex tuple. Hence this case is equivalent

By the mathematical induction, the lemma is proved.

III. THE PANPOSITIONABILITY OF Q_n^k , WHERE $k \geq 3$ IS AN odd integer and $n \geq 2$ is an integer.

Lemma 3. Q_2^3 is a panpositionable hamiltonian graph.

Proof: There are two cases: Case 1. u = (0,0) and v =(1,0); Case 2. u=(1,0) and v=(0,1). By brute force, we construct the required hamiltonian cycles. Please see Fig. 9.

Theorem 1. Q_n^3 is a panpositionable hamiltonian graph.

Proof: The theorem is proved by mathematical induction using Lemma 3 as base case. The detailed derivation is skipped.

Lemma 4. Let k be an odd integer with $k \geq 5$. Then Q_2^k is nearly-panpositionable.

Proof: The proof is by brute force and hence is skipped.

Theorem 2. Let k be an odd integer with $k \geq 5$. Q_n^k is nearlypanpositionable hamiltonian.

Proof: We will prove the theorem using the mathematical induction. By Lemma 4, Q_2^k is nearly-panpositionable hamiltonian. With the induction hypothesis, we assume that Q_{n-1}^k is nearly-panpositionable hamiltonian for some $n\geq 3.$ We need to show that Q_n^k is nearly-panpositionable hamiltonian. Let $u,v\in Q_n^k$ and l be an integer with

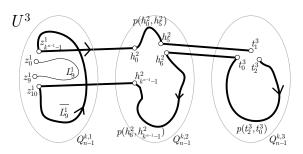


Fig. 10. U^3 for r=3 and l'=9.

 $d \leq l \leq \frac{|Q_n^k|}{2}$, where $d = d_{Q_n^k}(u, v)$.

Case 1. $u, v \in Q_{n-1}^{k,i}$. W.L.O.G., let i = 0.

Obviously, $d_{Q_{n-1}^k}(u,v) = d$.

Case 1.1. $d \le l \le \frac{k^{n-1}-1}{2}$. By the induction hypothesis, there exist a hamiltonian cycle $\begin{array}{l} C_l^i = \langle x_0^i, x_1^i, ..., x_l^i, ..., x_{k^{n-1}-1}^i, x_0^i \rangle \text{ in } Q_{n-1}^{k,i} \text{ for } u = x_0^0 \\ \text{and } v = x_l^0. \text{ Then we have the hamiltonian cycle } C = \\ \langle x_0^0, p(x_0^0, x_l^0), x_l^1, x_l^2, ..., x_l^{k-1}, p(x_{l+1}^{k-1}, x_{l-1}^{k-1}), \operatorname{inv}(p(x_{l+1}^{k-2}, x_{l-1}^{k-2})), p(x_{l+1}^{k-3}, x_{l-1}^{k-3}), \operatorname{inv}(p(x_{l+1}^{k-4}, x_{l-1}^{k-4})), ..., p(x_{l+1}^2, x_{l-1}^2), \\ \operatorname{inv}(p(x_{l+1}^1, x_{l-1}^1)), p(x_{l+1}^0, x_{n-1-1}^0), x_0^0 \rangle. \end{array}$

Case 1.2. $\frac{k^{n-1}-1}{2}+1 \leq l \leq \frac{|Q_k^k|}{2}$.

By the induction hypothesis, for any two vertices $x,y\in V(Q_{n-1}^k)$ and $1\leq l'\leq k^{n-1}-1$ there exists a hamiltonian cycle C of $Q_{n-1}^{k,i}$ with $d_C(x,y)=l'$. We set $x=z_0^i$ and $y=z_{l'}^i$, then the hamiltonian cycle will be $\langle z_0^i, p(z_0^i, z_{k^{n-1}-1}^i), z_0^i \rangle$. Split the hamiltonian cycle into two pathes $L^i_{l'}$ and $\overline{L^i_{l'}}$ by letting $L_{l'}^i = p(z_0^i, z_{l'}^i)$ and $L_{l'}^i = p(z_{l'+1}^i, z_{k^{n-1}-1}^i)$.

In [12], it is shown that for all $x, y \in Q_{n-1}^{k,i}$, there exists a hamiltonian path H^i of $Q_{n-1}^{k,i}$ between x and y. Define $H^i = p(h_0^i, h_{k^{n-1}-1}^i)$ with $x = h_0^i$ and $y = h_{k^{n-1}-1}^i$. By Lemma 2, for any path with length 2 denoted by $\langle t_0^i, t_1^i, t_2^i \rangle$, there exists a hamiltonian cycle $T^i = \langle t_0^i, p(t_0^i, t_{k^{n-1}-1}^i), t_0^i \rangle$ of $Q_{n-1}^{k,i}$. Let $t_0^i=h_{j+1}^i,\ t_1^i=h_j^i,\ z_{l'+1}^i=h_{k^{n-1}-1}^i, t_0^i,\ h_0^i=z_{k^{n-1}-1}^i$ and . Then there is a unique path $U^i=\langle t_2^i,p(t_2^i,t_0^i),h_{j+1}^{i-1},p(h_{j+1}^{i-1},h_{k^{n-1}-1}^{i-1}),z_{l'+1}^{i-2},\overline{L_{l'}^{i-2}},z_{k^{n-1}-1}^{i-2},\ h_0^{i-1},p(h_0^{i-1},h_j^{i-1}),t_1^i\rangle.$ For example, let r=3 and l'=9, then $U^3 = \langle t_2^3, p(t_2^3, t_0^3), h_6^2, p(h_6^2, h_{k^{n-1}-1}^2), z_{10}^1, L_9^1,$ $z_{k^{n-1}-1}^1, h_0^2, p(h_0^2, h_5^2), t_1^3 \rangle$. Please see Fig. 10 for an illustra-

Let m and r be integers and $0 \le r \le \frac{k-1}{2}$ such that $\frac{k^{n-1}-1}{2}-m+r\cdot k^{n-1}+l'+1=l.$ W.L.O.G., let $u=x_0^0$ and $v=x_{\frac{k^{n-1}-1}{2}-m}^{0}$. For simplicity, denote $x_{\frac{k^{n-1}-1}{2}-m}^{i}$ as $\begin{array}{c} v^i,\, x^i_{\frac{k^{n-1}-1}{2}-m-1} \text{ as } v^i_0 \text{ and } x^i_{\frac{k^{n-1}-1}{2}-m+1} \text{ as } v^i_1. \text{ If } r \text{ is even,} \\ \text{let } t^i_1 = v^i_0,\, t^i_2 = v^i,\, z^{1+r}_0 = x^{1+r}_0 \text{ and } z^{1+r}_{l'} = x^{1+r}_1. \text{ There is } \end{array}$

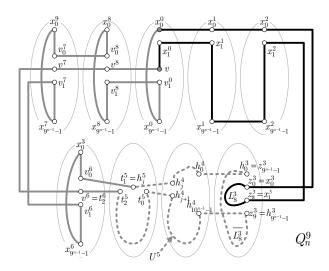


Fig. 11. An Example of Case 1.2 with k=9 and $l=\frac{5\cdot 9^{n-1}-17}{2}$.

a hamiltonian cycle

$$\begin{split} C &= & \langle x_0^0, x_0^1, ..., x_0^r, L_{l'}^{r+1}, p(x_1^r, x_{k^{n-1}-1}^r), \operatorname{inv}(p(x_1^{r-1}, x_{k^{n-1}-1}^{r-1})), p(x_1^{r-2}, x_{k^{n-1}-1}^{r-2}), \operatorname{inv}(p(x_1^{r-3}, x_{k^{n-1}-1}^{r-3})), \\ & ..., p(x_1^2, x_{k^{n-1}-1}^2), \operatorname{inv}(p(x_1^1, x_{k^{n-1}-1}^1)), p(x_0^1, v_0^0), v, \\ & v^{k-1}, v^{k-2}, ..., v^{r+4}, U^{r+3}, v_0^{r+4}, \operatorname{inv}(p(v_1^{r+4}, v_0^{r+4})), \\ & p(v_1^{r+5}, v_0^{r+5}), \operatorname{inv}(p(v_1^{r+6}, v_0^{r+6})), p(v_1^{r+7}, v_0^{r+7}), ..., \\ & p(v_1^{k-2}, v_0^{k-2}), \operatorname{inv}(p(v_1^{k-1}, v_0^{k-1})), v_1^0, p(x_{\frac{k^{n-1}-1}{2}-m+2}^0, x_{k^{n-1}-1}^0), x_0^0 \rangle. \end{split}$$

Please see Fig. 11 for an illustration, where m = 0, r = 2, l' = 8 and the hamiltonian cycle is

$$\begin{split} C &=& \langle x_0^0, x_0^1, x_0^2, L_8^3, p(x_1^2, x_{9^{n-1}-1}^2), \operatorname{inv}(p(x_1^1, x_{9^{n-1}-1}^1)), \\ & & p(x_1^0, v_0^0), v, v^8, v^7, v^6, U^5, v_0^6, \operatorname{inv}(p(v_1^6, v_0^6)), \\ & & p(v_1^7, v_0^7), \operatorname{inv}(p(v_1^8, v_0^8)), v_1^0, p(x_{\frac{9^{n-1}-1}{2}+2}^0, x_{9^{n-1}-1}^0), \\ & & x_0^0 \rangle. \end{split}$$

If r is odd, let $t_1^i = v_1^i$, $t_2^i = v^i$, $x_0^{1+r} = z_0^{1+r}$ and $x_{k^{n-1}-1}^{1+r} = z_0^{1+r}$ z_{ν}^{1+r} . There is a hamiltonian cycle

$$\begin{split} C &=& \langle x_0^0, x_0^1, ..., x_0^r, L_{\mathcal{V}}^{r+1}, \mathrm{inv}(p(x_1^r, x_{k^{n-1}-1}^r)), p(x_1^{r-1}, \\ & x_{k^{n-1}-1}^{r-1}), \mathrm{inv}(p(x_1^{r-2}, x_{k^{n-1}-1}^{r-2})), p(x_1^{r-3}, x_{k^{n-1}-1}^{r-3}), \\ & ..., \mathrm{inv}(p(x_1^3, x_{k^{n-1}-1}^3)), p(x_1^2, x_{k^{n-1}-1}^2), \mathrm{inv}(p(x_1^1, x_{k^{n-1}-1}^1)), p(x_1^0, v_0^0), v, v^{k^{n-1}-1}, v^{k^{n-1}-2}, ..., v^{r+4}, \\ & U^{r+3}, v_0^{r+4}, p(v_1^{r+4}, v_0^{r+4}), \mathrm{inv}(p(v_1^{r+5}, v_0^{r+5})), \\ & p(v_1^{r+6}, v_0^{r+6}), \mathrm{inv}(p(v_1^{r+7}, v_0^{r+7})), ..., \\ & p(v_1^{k^{n-1}-2}, v_0^{k^{n-1}-2}), \mathrm{inv}(p(v_1^{k^{n-1}-1}, v_0^{k^{n-1}-1})), \\ & v_1^0, p(x_{k^{n-1}-1}^0, x_0^{n-1}, x_0^{n-1}), x_0^0 \rangle. \end{split}$$

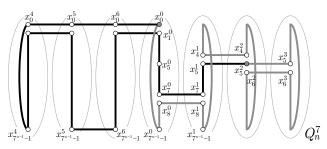


Fig. 12. An Example of Case 2.1 with k = 7 and $l = 3 \cdot 7^{n-1} + 11$.

of the shortest path between x_0^0 and $v=x_{d'}^0$ in $Q_{n-1}^{k,0}$. Note that for all $i\leq k-1$, there is a hamiltonian cycle $C^i=\langle x_0^i,x_1^i,...,x_{d'}^i,...,x_{k^{n-1}}^i\rangle$ in $Q_{n-1}^{k,i}$.

Case 2.1. l-d is even.

Let $0 \le r \le \frac{k-1}{2}$ be an integer, $d+2(t-d')+r\cdot k^{n-1}=l$, $d' \le t \le k^{n-1}-1$ and e=k-1-r.

Let r be an odd integer. We have the hamiltonian cycle $C = \langle x_0^0, x_0^{k-1}, x_0^{k-2}, ..., x_0^{k-r}, \operatorname{inv}(p(x_1^{k-r}, x_{k^{n-1}-1}^{k-r})), p(x_1^{k-r+1}, x_{k^{n-1}-1}^{k-r+1}), \operatorname{inv}(p(x_1^{k-r}, x_{k^{n-1}-1}^{k-r+1})), p(x_1^{k-r+1}, x_{k^{n-1}-1}^{k-r+1}), \operatorname{inv}(p(x_1^{k-3}, x_{k^{n-1}-1}^{k-r+3})), p(x_1^{k-2}, x_{k^{n-1}-1}^{k-r}), \operatorname{inv}(p(x_1^{k-3}, x_{k^{n-1}-1}^{k-r+3})), p(x_1^{k-2}, x_{k^{n-1}-1}^{k-r}), \operatorname{inv}(p(x_1^{k-1}, x_{k^{n-1}-1}^{l-r})), p(x_1^{k-2}, x_{k^{n-1}-1}^{k-r}), \operatorname{inv}(p(x_1^{k-1}, x_{d^r}^{l-r})), p(x_1^{k-1}, x_{d^r-1}^{l-r})), \operatorname{inv}(p(x_{d^r+1}^{l-r}, x_{d^r-1}^{l-r})), p(x_{d^r+1}^{l-r}, x_{d^r-1}^{l-r})), \operatorname{inv}(p(x_{d^r+1}^{l-r}, x_{d^r-1}^{l-r})), p(x_{d^r+1}^{l-r}, x_{d^r-1}^{l-r})), p(x_{d^r+1}^{l-r}, x_{d^r-1}^{0}), p(x_{d^r+1}^{l-r}, x_{d^r-1}^{l-r})), p(x_{d^r+1}^{l-r}, x_{d^r-1}^{0}), p(x_{d^r+1}^{l-r}, x_{d^r-1}^{l-r})), p(x_{d^r+1}^{l-r}, x_{d^r-1}^{0}), p(x_{d^r+1}^{l-r}, x_{d^r-1}^{l-r})), p(x_{d^r+1}^{l-r}, x_{d^r-1}^{l-r}), p(x_{d^r+1}^{l-r}, x_{d^r-1}^{l-r})), p(x_{d^r+1}^{l-r}, x_{d^r-1}^{l-r}), \operatorname{inv}(p(x_{d^r+1}^{l-r}, x_{d^r-1}^{l-r})), p(x_{d^r+1}^{l-r}, x_{d^r-1}^{l-r})), p(x_{d^r+1}^{l-r}, x_{d^r+1}^{l-r}), \operatorname{inv}(p(x_{d^r+1}^{l-r}, x_{d^r-1}^{l-r})), p(x_{d^r+1}^{l-r}, x_{d^r-1}^{l-r}), \operatorname{inv}(p(x_{d^r+1}^{l-r}, x_{d^r-1}^{l-r})), p(x_{d^r+1}^{l-r}, x_{d^r-1}^{l$

 $\begin{array}{l} p(x_{6}^{-},x_{4}^{-}), \operatorname{inv}(p(x_{8}^{-},x_{4}^{-})), p(x_{8}^{-},x_{0}^{-}), x_{0}^{-}). \\ \operatorname{Let} r \text{ be an even integer. We have the amiltonian cycle } C = \\ \langle x_{0}^{0}, x_{0}^{k-1}, x_{0}^{k-2}, \ldots, x_{0}^{k-r}, p(x_{1}^{k-r}, x_{k^{n-1}-1}^{k-r}), \operatorname{inv}(p(x_{1}^{k-r+1}, x_{k^{n-1}-1}^{k-r+1})), p(x_{1}^{k-r+2}, x_{k^{n-1}-1}^{k-r+1}), \operatorname{inv}(p(x_{1}^{k-r}, x_{k^{n-1}-1}^{k-r+1})), \ldots, \\ p(x_{1}^{k-2}, x_{k^{n-1}-1}^{k-2}), \operatorname{inv}(p(x_{1}^{k-1}, x_{k^{n-1}-1}^{-1})), x_{1}^{0}, p(x_{2}^{0}, x_{d'-1}^{0}), \\ p(x_{0}^{0}, x_{1}^{0}), \operatorname{inv}(p(x_{d}^{1}, x_{1}^{1})), x_{d}^{2}, \ldots, x_{d'}^{2}, \ldots, x_{d'}^{e}, p(x_{d'+1}^{e}, x_{d'-1}^{e}), \\ \operatorname{inv}(p(x_{d'+1}^{e-1}, x_{d'-1}^{e-1})), p(x_{d'+1}^{e-2}, x_{d'-1}^{e-2}), \operatorname{inv}(p(x_{d'+1}^{e-3}, x_{d'-1}^{e-3})), \ldots, \\ p(x_{d'+1}^{2}, x_{d'-1}^{2}), \operatorname{inv}(p(x_{1+1}^{1}, x_{d'-1}^{1})), p(x_{1}^{0}, x_{0}^{0}). \end{array}$

Case 2.2. l - d is odd.

By the induction hypothesis, there exists a hamiltonian cycle $D^i=\langle y_0^i,y_1^i,...,y_{k^{n-1}}^i\rangle$ in $Q_{n-1}^{k,i}$ such that $x_0^i=y_0^i,x_{d'}^i=y_{l'}^i$ and the length l' of the path joining y_0^i to $y_{l'}^i$ in $Q_{n-1}^{k,i}$ is the smallest integer when l'-d' is odd. Let $0\leq r\leq \frac{k-1}{2}$ be an integer, $d+l'-d'+2(t-l')+r\cdot k^{n-1}=l,$ $d'\leq t\leq k^{n-1}-1$ and e=k-1-r.

and c=n-1 . Let r be an odd integer. We have the hamiltonian cycle $C=\langle y_0^0, y_0^{k-1}, y_0^{k-2}, ..., y_0^{k-r}, \operatorname{inv}(p(y_1^{k-r}, y_{k^{n-1}-1}^{k-r})), p(y_1^{k-r+1}, y_{k^{n-1}-1}^{k-r+1}), \operatorname{inv}(p(y_1^{k-r+2}, y_{k^{n-1}-1}^{k-r+2})), p(y_1^{k-r+3}, y_{k^{n-1}-1}^{k-r+3}), ..., [5] \\ \operatorname{inv}(p(y_1^{k-3}, y_{k^{n-3}-1}^{k-3})), p(y_1^{k-2}, y_{k^{n-1}-1}^{k-r}), \operatorname{inv}(p(y_1^{k-1}, y_{k^{n-1}-1}^{k-r})), \\ y_1^0, p(y_2^0, y_{l'-1}^0), p(y_{l'}^0, y_0^1), \operatorname{inv}(p(y_{l'}^1, y_l^1)), y_{l'}^2, ..., y_{l'}^j, ..., y_{l'}^e, \\ \operatorname{inv}(p(y_{l'+1}^e, y_{l'-1}^e)), p(y_{l'+1}^e, y_{l'-1}^e), \operatorname{inv}(p(y_{l'+1}^e, y_{l'-1}^e)), \\ p(y_{l'+1}^e, y_{l'-1}^e), ..., \operatorname{inv}(p(y_{l'+1}^1, y_{l'-1}^1)), p(y_{l'+1}^0, y_{l'}^0), y_0^0).$

Let r be an even integer. We have the hamiltonian cycle $C=\langle y_0^0,y_0^{k-1},y_0^{k-2},...,y_0^{k-r},p(y_1^{k-r},y_{k^{n-1}-1}^{k-r}),\operatorname{inv}(p(y_1^{k-r+1},y_{k^{n-1}-1}^{k-r+1})),p(y_1^{k-r+2},y_{k^{n-1}-1}^{k-r+2}),\operatorname{inv}(p(y_1^{k-r+3},y_{k^{n-1}-1}^{k-r+3})),...,p(y_1^{k-2},y_{k^{n-1}-1}^{k-2}),\operatorname{inv}(p(y_1^{k-1},y_{k^{n-1}-1}^{-1})),y_1^0,p(y_2^0,y_{l'-1}^0),p(y_{l'}^0,y_0^1,\operatorname{inv}(p(y_{l'}^1,y_{l'}^1)),y_{l'}^2,...,y_{l'}^j,...,y_{l'}^e,p(y_{l'+1}^e,y_{l'-1}^e),\operatorname{inv}(p(y_{l'+1}^e,y_{l'-1}^e)),p(y_{l'+1}^e,y_{l'-1}^e),\operatorname{inv}(p(y_{l'+1}^e,y_{l'-1}^e),\operatorname{inv}(p(y_{l'+1}^1,y_{l'-1}^e)),p(y_{l'+1}^0,y_{l'-1}^e),p(y_{l'+1}^0,y_{l'}^0,y_0^0),0).$

By the mathematical induction, the theorem is proved.

IV. Q_n^k is bipanpositionable, where $k \geq 4$ is an even integer and $n \geq 2$ is an integer.

Lemma 5. Let k be an even integer with $k \geq 4$. Then Q_2^k is bipanpositionable.

Proof: The proof is by brute force and hence is skipped.

Theorem 3. Let k be an even integer with $k \geq 4$. Q_n^k is bipanpositionable hamiltonian.

Proof: We will prove the theorem using the mathematical induction. By Lemma 5, Q_2^k is bipanpositionable hamiltonian. With the induction hypothesis, we assume that Q_{n-1}^k is bipanpositionable hamiltonian for some $n \geq 4$. We need to show that Q_n^k is bipanpositionable hamiltonian. Note that Q_n^k is bipanpositionable [7]. Therefore, $Q_n^4 = Q_{2n}^2$ is bipanpositionable. It suffices to prove the cases for $k \geq 6$.

The proof is similar to Theorem 2, so we'll skip it. Readers can follow the similar techniques in Theorem 2 to construct the required hamiltonian cycles.

V. CONCLUSIONS

In this paper, we prove that the k-ary n-cube Q_n^3 is panpositionable hamiltonian and Q_n^k is nearly-papositionable for any odd integer $k \geq 5$. Moreover, we prove that Q_n^k is bipanpositionable hamiltonian for any even integer $k \geq 4$. It is known that the hypercube, Q_n^2 , is bipanpositionable [7]. Thus Q_n^k is bipanpositionable for all even integers $k \geq 2$.

The panpositionability of any k-ary n-cube has been completely studied and the result is optimal in the sense that given any two vertices u and v, there exists no more hamiltonian cycle on which $d_C(u,v)$ equals any of the numbers we miss in the nearly-panpositionable Q_n^k when k is odd.

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