Characterizations of Star-Shaped, L-Convex, and Convex Polygons

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Abstract—A chord of a simple polygon $P$ is a line segment $[xy]$ that intersects the boundary of $P$ only at both endpoints $x$ and $y$. A chord of $P$ is called an interior chord provided the interior of $[xy]$ lies in the interior of $P$. $P$ is weakly visible from $[xy]$ if for every point $v$ in $P$ there exists a point $w$ in $[xy]$ such that $[vw]$ lies in $P$. In this paper star-shaped, L-convex, and convex polygons are characterized in terms of weak visibility properties from internal chords and star-shaped subsets of $P$. A new Krasnoselskii-type characterization of isothetic star-shaped polygons is also presented.

Keywords—Convex polygons, L-convex polygons, star-shaped polygons, chords, weak visibility, discrete and computational geometry.

I. INTRODUCTION

THIS paper is concerned mainly with new characterizations of star-shaped and L-convex polygons. However, as a corollary a new characterization of convex polygons is also obtained. For any integer $n \geq 3$, a polygon in the Euclidean plane $E^2$ is defined as the figure $P = [x_1, x_2, ..., x_n]$ formed by $n$ points $x_1, x_2, ..., x_n$ in $E^2$ and $n$ line segments $[x_i, x_{i+1}]$, $i = 1, 2, ..., n-1$, and $[x_n x_1]$. The points $x_i$ are called the vertices of the polygon and the line segments are termed its edges. The vertices of $P$ are assumed to be in general position, i.e., no three vertices are collinear. For an accessible introduction to polygons and their classification see the paper by Grunbaum [13].

Definition: A polygon $P$ is called a simple polygon provided that no point of the plane belongs to more than two edges of $P$ and the only points of the plane that belong to precisely two edges are the vertices of $P$.

A simple polygon $P$ has a well defined boundary denoted by $bd(P)$, an interior denoted by $int(P)$, and an exterior denoted by $ext(P)$. By convention, the interior of a polygon is provided the interior of $[xy]$ lies in $P$. The vertices of a simple polygon are of two types: convex and concave. In the mathematics literature the terminologies reentrant vertex and local non-convexity point are often used instead of concave vertex, whereas in the computational geometry literature the word reflex vertex is preferred. However, in this paper the more natural term concave is used. For a given vertex $x_j$ let $y = \lambda x_{j-1} + (1-\lambda)x_j$ and $z = \mu x_{j-1} + (1-\mu)x_j$. For all sufficiently small positive values of $\mu$ and $\lambda$ $int([xz])$ lies either totally in $int(P)$ or totally in $ext(P)$; in the former case $x_j$ is a convex vertex whereas in the latter case it is a concave vertex.

Definition: A simple polygon $P$ is called star-shaped if there exists a point $x \in P$ such that for all points $y \in P$, $[xy] \subseteq P$. The collection of all such points $x$ is called the kernel of $P$.

Definition: (Grunbaum [13]) A simple polygon $P$ is called convex provided that all its vertices are convex.

This is a special case of a well-known theorem due to Tietze [29], which states that if $S$ is a closed connected set in a Euclidean space, all of whose points are points of local convexity, then $S$ is convex. A point $x \in S$ is a point of local convexity of $S$ if there exists a neighborhood $N$ of $x$ such that $N \cap S$ is convex; otherwise $x$ is called a point of local non-convexity of $S$.

The convexity of polygons is frequently characterized in terms of the connectivity of the intersection-sets of lines that intersect the polygon Fary [11]. More recently Pinelis characterized the convexity of cyclic polygons in terms of the central angles of the polygon [23]. Closer in spirit to the results presented here, Nagel characterizes convex polygons in terms of the orientation-dependent chord length distributions [22].

Characterizations of objects such as convex and star-shaped polygons as well as more general sets are of interest for at least two reasons. Mathematicians are motivated by the desire to obtain a deeper understanding of geometric objects such as polygons. Different characterizations of an object provide different views of the object and thus further this understanding [11], [16], [27], [35], [36]. The study of convex sets is relevant to a variety of disciplines in science and technology [17]. For example, computer scientists are interested in designing algorithms for recognizing whether polygons are convex, L-convex, star-shaped, etc., in a variety of contexts driven by applications in pattern recognition and computer vision problems. Different characterizations yield alternative algorithms, with varying computational complexities, for solving such problems [2, 3, [12], [25], [30].

A simple polygon $P$ is also said to be convex if every pair of points $x, y \in P$ can be joined by a line segment $[xy] \subseteq P$ [11]. This well known characterization of convex polygons is equivalent to the demand that all three of the segments determined by each triplet of pairwise distinct points in $P$ lie totally in $P$. One can relax this criterion and still obtain a characterization of convex polygons. A simple polygon $P$ is convex if, and only if, it contains two of the three segments determined by each triplet of its points [20]. Further weakening the criterion to the new demand that only one of the three segments be contained in $P$ does not lead to convexity but to

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the notion of $P_3$-convexity [32]. Valentine [32] has shown that a $P_3$-convex polygon can be represented as the union of three or fewer convex polygons.

Convex polygons have also been characterized in terms of nearest point properties [35], as illustrated by the Theorem of Bunt-Motzkin [37] which states that a simple polygon $P$ is convex if, and only if, for every point $p$ in $P$ there is exactly one point of $P$ nearest to $p$.

Furthermore, there has been interest in characterizing convex polygons in terms of unimodality properties. There are several possibilities for definitions of the notion of unimodality depending on the distance functions employed. For example, one can define for a vertex $z$ of $P$, a function $f(z)$ which is the Euclidean distance between $z$ and each vertex of $P$ in the order in which the vertices occur in $P$. If $f(z)$ is unimodal then $z$ is called a unimodal vertex. It has been incorrectly assumed in several published papers that a polygon is convex if all its vertices are unimodal in this sense. Furthermore algorithms for computing geometric properties of convex polygons based on this assumption have also been published. However, counter examples to the claim [2] and to such algorithms [3] have since appeared. Just as the Euclidean distance between pairs of vertices is used to create $f(z)$ one may instead consider vertex-edge or edge-vertex pairs and measure the separation by means of the perpendicular distance between the vertex and the line collinear with the edge in question. In this way for an edge $e$ of $P$ one can define a function $g(e)$ which is the perpendicular distance from the line collinear with $e$ to every vertex of $P$ in the order in which the vertices appear in $P$. If $g(e)$ is unimodal then $e$ called a unimodal edge. In [30] it is shown that if all the edges of $P$ are unimodal in this sense then $P$ is convex.

A type of unimodality that is very different to that discussed above was considered by Dharmadhikari and Jogdeo [8]. Let $P$ be a simple polygon in $R^2$. Given a non-zero vector $u \in R^2$ and $k \in R$, denote by $L(u, k)$ the line determined by the dot product $u \cdot x = k$. Let $f_u(k)$ denote the measure of $P \cap L(u, k)$. If $P$ is a convex polygon then $f_u(k)$ as a function of $k$, is first non-decreasing and then non-increasing for every value of $u$. A non-negative function $f$ on $R$ is said to be unimodal if there exits a $v \in R$ such that $f$ is non-decreasing on $(-\infty, v]$ and non-increasing on $[v, \infty)$. Furthermore such a number $v$ need not be unique. Consider now the following condition:

**Condition A:** For every fixed non-zero $u \in R^2$, the function $f_u(k)$ is unimodal in $k$.

It is natural to ask whether condition $A$ is sufficient for a simple polygon $P$ to be convex. The answer to this question is negative. Consider the following example from [8], and refer to Fig. 1. Let $ABCD$ and $A'B'C'D'$ be parallelograms which are mirror images of each other and are such that the lines through $CD$ and $C'D'$ meet outside the polygon at some point $z$. It can easily be verified by inspection that this polygon satisfies condition $A$ but it is not even star-shaped; no point in $P$ sees both $C$ and $C'$. On the other hand in [8] it is shown that if $P$ is such that for every fixed $u \in R^2$, the function $f_u(k)$ is continuous on the interior of its support then $P$ is convex.

The earliest characterization of star-shaped polygons is due to Krasnoselskii [18], and this result is known as Krasnoselskii’s theorem [37]. This theorem states that if for every set of three points $x, y$, and $z \in P$ there exists a point $w \in P$ (possibly dependent on $x, y$, and $z$), such that the three segments $[wx], [wy], [wz]$ all lie in $P$ then $P$ is star-shaped. If a polygon is *isothetic* then it has been shown that the number three can be reduced to two [31]. A polygon $P$ is called *isothetic* provided that all its edges are parallel to either the $x$ or the $y$ axis. For additional Krasnoselskii-type characterizations of star-shaped polygons the reader is referred to [4], [21], [24], and [33]. A different characterization is provided in [26], where it is shown that a polygon $P$ is star-shaped if, and only if, the intersection of all the maximal convex sub-polygons of $P$ is non-empty.

### II. Weak Visibility Characterizations

In this section new characterizations of star-shaped and convex polygons are presented based on the notion of weak visibility. A polygon $P$ is said to be weakly visible [1] from a subset $S$ of $P$ if for every point $x \in P$ there exists a point $y \in S$ such that the line segment $[xy] \in P$. A chord of a polygon $P$ is a line segment $[xy]$ that intersects the boundary $bd(P)$ only at $x$ and $y$. If the interior of $[xy]$ lies in the interior of $P$ then the chord is termed an *interior chord*. If the interior of $[xy]$ lies in the exterior of $P$ then the chord is said to be an *exterior chord*.

**Theorem 2.1:** A simple polygon $P$ is star-shaped if, and only if, there exists a point $x \in P$ such that $P$ is weakly visible from every internal chord traversing $x$.

**Proof:** (only if part) Choose $x$ to be any point in the kernel of $P$. Since $P$ is star-shaped from $x$ it follows that it is weakly visible from every internal chord of $P$ that traverses $x$.

(If part) Assume $P$ contains a point $x$ such that $P$ is weakly visible from every internal chord traversing $x$. Then $P$ must be star-shaped from $x$. If this were not so it would imply the existence of a point $y$ in $P$ that is not visible from $x$. Now construct a line through both $x$ and $y$ and let $a, b$ denote the first points of intersection of $L$ with $bd(P)$ as $L$ is traversed in both directions starting at $x$. Let $L' \in L$ denote the segment $[ab]$. Since $y, a$, and $b$ all lie on $L$ and $y$ is not visible from $x$ it follows that $y$ is not visible from any point on $L'$. Therefore $P$ is not weakly visible from chord $L'$ which is a contradiction.

![Fig. 1 Condition A is not sufficient to ensure either convexity](image)
Note that as a corollary a new characterization of convex polygons is obtained. The theorem actually proves that if a point \( x \) exists such that \( P \) is weakly visible from every internal chord traversing \( x \) then \( P \) is star-shaped from \( x \). Thus if this property holds true for every point \( x \in P \), it follows that \( P \) is star-shaped from every point in \( P \), and is convex. Thus the following result follows.

**Corollary 2.1:** A simple polygon \( P \) is convex if, and only if, it is weakly visible from every internal chord of \( P \).

Consider now the case of isothetic polygons, i.e., polygons with all their edges parallel to the coordinate axes. Such a polygon with its four types of tabs is illustrated in Fig. 2.

![An isothetic polygon and its four types of tabs](image)

A **tab** is a set of two adjacent convex vertices along with the three edges of \( P \) incident on these two vertices. There are four types of tabs. For example, in Fig. 2, \([a, b, c, d]\) is a **top tab**.

Recall that Theorem 2.1 states that an arbitrary simple polygon \( P \) is star-shaped if, and only if, there exists a point \( x \in P \) such that \( P \) is weakly visible from every internal chord traversing \( x \). The term **unoriented direction** \( q \) refers to an equivalence class of parallel lines that make an angle of \( q \) with respect to some agreed upon fixed axis. Also observe that in an arbitrary simple polygon \( P \) each of its edges can occur in any one of an infinite number of un-oriented directions. Return now to the case of isothetic polygons. It is natural to conjecture the following result analogous to theorem 2.1.

**Conjecture:** A simple isothetic polygon \( P \) is star-shaped if, and only if, there exists a point \( x \in P \) such that \( P \) is weakly visible from both the horizontal and vertical internal chords traversing \( x \). As it turns out however one can prove a stronger result for isothetic polygons in the form of Theorem 2.2.

**Theorem 2.2:** A simple isothetic polygon \( P \) is star-shaped if, and only if, \( P \) is weakly visible from **both some horizontal** and some **vertical internal chord** of \( P \).

**Proof:** (only if part) Let \( x \) be a point in the kernel of \( P \). Clearly \( P \) is weakly visible from any internal chord traversing point \( x \). Therefore \( P \) is weakly visible from both some horizontal and some vertical internal chord of \( P \), namely the horizontal and vertical chords traversing point \( x \).

Q.E.D.

The result follows immediately by recalling that a simple polygon \( P \) is weakly visible from every internal chord of \( P \) if, and only if, \( P \) contains vertices not visible from \([tb]\) where \( t \) and \( b \) are the upper and lower endpoints, respectively, of the chord. It follows that \( P \) must contain only one top tab and one bottom tab and that \([tb]\), the horizontal chord from which \( P \) is weakly visible must have its end-points \( l, r \) on the unique left and right tabs, respectively. Furthermore, the right tab of \( P \) must lie to the right of \([tb]\) and the left tab of \( P \) must lie to the left of \([tb]\). For assume this not to be the case and, without loss of generality, let \( P \) contain a right tab to the left of \([tb]\). This would imply that \( P_1 \) or \( P_2 \) contains a top or bottom tab other than those determined by \([tb]\) which in turn would contradict the fact that \( P \) is weakly visible from \([tb]\). Therefore every internal chord that weakly sees \( P \) will intersect each such vertical chord. It still remains to show that \( P \) is star-shaped. Let \( z \) be the intersection point of \([tb]\) and \([lr]\). It will be shown that \( P \) is star-shaped from \( z \). Assume \( P \) is not star-shaped from \( z \). This implies that there exists a point \( w \) on bd(\( P \)) that is not visible from \( z \). Without loss of generality assume that \( w \) lies on that part of bd(\( P \)) between \( r \) and \( b \) as \( P \) is traversed in a clockwise manner, and denote this portion of \( P \) by Chain[\( r, ..., b \)]. Let \( P_{arb} = Chain[r, ..., b] \cup [b, z] \cup [z, r] \) and let VP(\( P_{arb}, z \)) be the visibility region of \( P_{arb} \) from \( z \). The region VP(\( P_{arb}, z \)) cuts off regions of \( P_{arb} \) which are hidden from \( z \) and lie either to the left or to the right of the cutting visibility rays emanating from \( z \). Clearly \( w \) must lie either in a left or a right such hidden region. In the former case \( w \) is not visible from \([lr]\), and in the latter case \( w \) is not visible from \([tb]\). In both cases a contradiction results, and therefore \( P \) must be star-shaped from \( z \).

Q.E.D.

It has been shown in Theorem 2.2 that the point \( x \) in the conjecture could be disposed of and that it was sufficient to impose weak visibility from some horizontal and some vertical chord in order to characterize isothetic star-shaped polygons. This opens a similar question for the original non-isothetic simple polygons. In other words, is it true that an arbitrary simple polygon \( P \) is star-shaped if, and only if, for every un-oriented direction \( q \) there exists an internal chord of \( P \) with direction \( q \) from which \( P \) is weakly visible? The answer to this question is negative and a counterexample due to ElGindy [10] is illustrated in Fig. 3. Apart from three thin spikes at \( a, b, \) and \( c \) the polygon in Fig. 3 has its remaining vertices on a circle with center \( z \). Furthermore the spikes are so thin and so placed that their extended visibility lines form a triangle that encloses \( z \). It is clear by observation that for every unoriented direction there exists an internal chord of \( P \) with direction \( q \) from which \( P \) is weakly visible. Consider any chord of \( P \) with direction \( q \), and passing through \( z \). If \( q \) is not...
one of the directions in the set determined by the three spikes, then \( P \) is weakly visible from this chord. On the other hand if \( q \) is contained in one such set, say that of spike \( a \), then translate the chord in a direction orthogonal to \( q \) until it intersects the visibility cone of spike \( a \).

**Theorem 3.2:** (Toussaint & ElGindy [31]) If every two points on the boundary of an *isothetic* simple polygon \( P \) are visible from some common point in \( P \) then there exists a point in \( P \) from which the entire boundary of \( P \) is visible.

In order to proceed straightforwardly some definitions for arbitrary simple polygons are introduced.

**Definition:** The two closed rays of a line \( L \) which have only a point \( x \in L \) in common are called complementary rays. If \( R(x) \) is a ray with endpoint \( x \), its complementary ray is denoted by \( R'(x) \).

**Definition:** A set of rays is said to be *concurrent* if there exists a point of the plane that intersects each and every ray in the set.

**Definition:** A ray \( R(x) \) with endpoint \( x \in bd(P) \) is an external ray of support to \( int(P) \) if \( R(x) \cap int(P) = \emptyset \).

**Definition:** If \( x \in bd(P) \) then \( K(x) \) is the union of all the external rays of support to \( int(P) \) at \( x \). The set \( K(x) \) is called an *external cone of support*. The union of all the complementary rays \( R(x) \) where \( R(x) \in K(x) \) is denoted by \( K(P) \).

Valentine [33] proved the following result.

**Theorem 3.3:** (Valentine [33]) Let \( P \) be a non-convex simple polygon. Suppose that for each set of *three* (not necessarily distinct) concave vertices \( x_1, x_2, x_3 \) of \( P \), there exist three external rays of support at \( x_1, x_2, x_3 \) respectively to \( int(P) \) whose corresponding complementary rays are concurrent and meet in \( P \). Then \( P \) is star-shaped.

The main result of this section will now be proved.

**Theorem 3.4:** Let \( P \) be a non-convex simple *isothetic* polygon. Suppose that for each set of *two* (not necessarily distinct) concave vertices \( x_1, x_2 \) of \( P \), there exist two external rays of support at \( x_1 \) and \( x_2 \), to \( int(P) \) whose corresponding complementary rays are concurrent and meet in \( P \). Then \( P \) is star-shaped.

**Proof:** Let \( x_1 \) and \( x_2 \) be two concave vertices of \( P \). By hypothesis there exist two external rays of support at \( x_1, x_2 \), respectively, to \( int(P) \) whose corresponding complementary rays are concurrent and meet in \( P \) at some point \( z \). By construction it follows that \( [x_1z] \in K'(x_1) \) and \( [x_2z] \in K'(x_2) \). Therefore \( z \in K'(x_1) \cap K'(x_2) \). Since \( P \) is isothetic \( K'(x_1) \) and \( K'(x_2) \) can each be expressed as the intersection of the half-planes (containing \( z \)) determined by \( bd(K'(x_1)) \) and \( bd(K'(x_2)) \). Denote the vertical and horizontal such half-planes by \( V(z_i) \) and \( H(z_i) \), respectively, where \( i = 1, 2 \). Clearly the pair \( V(z_1) \) and \( V(z_2) \) must contain a non-zero intersection. Since the above arguments are true for all pairs of concave vertices of \( P \) it follows that all pairs of vertical half-planes intersect. From Helly’s theorem [14] it follows that all such vertical half-planes contain a non-zero intersection. Similar arguments hold for the horizontal half-planes. Therefore the intersection of \( K'(x_1) \) over all concave vertices \( x_1 \) must be non-zero. Denote this intersection by \( K^* \). Next it is shown that the intersection of \( K^* \) with \( P \) is non-empty. First note that \( P \) cannot contain more than one of each of the four types of tabs. For if this were not so, one type would contain at least two tabs and this would imply that there exist at least two concave vertices with external rays of support whose corresponding complementary

![Fig. 3 ElGindy's counterexample](image-url)

![Fig. 4 A non-star-shaped polygon with chords in the three directions from which the polygon is weakly visible](image-url)

### III. Krasnoselskii-Type Characterizations

In this section two characterizations of isothetic star-shaped polygons are presented that resemble Krasnoselskii’s theorem [18] for arbitrary simple polygons in their combinatorial flavor. First the original theorems for arbitrary and isothetic polygons are stated.

**Theorem 3.1:** (Krasnoselskii [18]) If every *three* points on the boundary of a simple polygon \( P \) are visible from some common point in \( P \) then there exists a point in \( P \) from which the entire boundary of \( P \) is visible.
rays are not concurrent, thus contradicting the hypothesis. Let \( K^{**} \) denote the intersection of \( K^* \) with the four interior half-planes determined by each of the four tabs. Each such half-plane contains \( P \) and is bounded by the line collinear with the two convex vertices making up the corresponding tab. It is a straightforward matter to show that if \( K^* \) is bounded, this intersection operation will not change \( K^* \), and if it is unbounded then \( K^{**} \) will be bounded but non-zero. Note also that the above discussion implies that each convex vertex of \( P \) not belonging to a tab must be such that both of its adjacent vertices are concave. This in turn implies that \( K^{**} \) is the intersection of the interior half-planes determined by all the edges of \( P \) and therefore \( P \) has a non-zero kernel. Therefore \( P \) is star-shaped. Q.E.D.

IV. \textit{L}-Convex Polygons

In 1949 Horn and Valentine introduced the definition of a \textit{link-distance} between two points \( a, b \) in \( P \) [15]. Since then mathematicians have investigated several properties of this distance measure [5, 6], [34], whereas computer scientists have investigated its computational aspects [19, 28]. The link-distance measure [5, 6], [34], whereas computer scientists have investigated its computational aspects [19, 28]. The link-distance measure [5, 6], [34], whereas computer scientists have investigated its computational aspects [19, 28].

A polygon \( P \) is said to be \( L \)-convex (or simply \( L \)-convex) if every pair of points \( a, b \) in \( P \) are \textit{link-distance} two apart. More generally \( P \) is said to be \( L \)-convex if every pair of points \( a, b \) in \( P \) have link-distance \( k \) between them. \( L \)-convex polygons have received some attention in the computational geometry literature. In particular, ElGindy, Avis and Toussaint [9] have shown that if a polygon is known to be \( L \)-convex it can be triangulated with a very simple algorithm in linear time. No such practical efficiency is known for arbitrary simple polygons, although a rather complicated linear-time algorithm was discovered by Chazelle [7], ElGindy, Avis and Toussaint [9] also showed that testing a simple polygon with \( n \) vertices for \( L \)-convexity can be done in \( O(n^3) \) time. Castiglioni et al. [6] investigate a special class of \( L \)-convex polygons called \( L \)-convex polyominoes from a tomographical point of view, and characterize them by means of horizontal and vertical projections. Horn and Valentine [15] proved that if \( P \) is \( L \)-convex then for every point \( x \) in \( P \) there exists a chord that traverses \( x \), say \( L(x) \), such that \( P \) is weakly visible from \( L(x) \). Since the converse also holds true this is in fact a characterization of \( L \)-convex polygons. An interesting question arises when one relaxes the chord \( L(x) \) traversing \( x \) to allow more general regions such as \textit{star-shaped} regions.

Horn and Valentine [15] characterized \( L \)-convex polygons in terms of a covering of \( P \) expressed by the following theorem.

**Theorem 4.1:** (Horn & Valentine [15]) A simple polygon \( P \) is \( L \)-convex if, and only if, \( P \) can be expressed as the sum of convex subsets of \( P \), every two of which have a point in common.

Here an alternate characterization of \( L \)-convex polygons is obtained in terms of weak visibility. In the sequel let \( S^*(x) \) denote a star-shaped subset of \( P \) containing \( x \) from which \( P \) is weakly visible.

**Theorem 4.2:** A simple polygon \( P \) is \( L \)-convex if, and only if, \( P \) has the property that for every point \( x \) in \( P \) there exists a subset \( S^* \) of \( P \) such that: (1) \( x \) is contained in \( S^* \), (2) \( S^* \) is star-shaped from \( x \), and (3) \( P \) is weakly-visible from \( S^* \).

**Proof:** [only if] If \( P \) is \( L \)-convex it has the property that for every point \( x \) in \( P \) there exists a traversing chord \( L(x) \) from which \( P \) is weakly visible [15]. Clearly \( L(x) \) satisfies the three conditions of the theorem.

[iif] Let \( x \) and \( y \) be any two points in \( P \). The weak visibility of \( P \) from \( S^*(x) \) implies that there exists a point \( z \) in \( S^*(x) \) visible from \( y \). From the star-shapedness of \( S^*(x) \) from \( x \) it follows that \( x \) and \( z \) are visible. Therefore \( x \) and \( y \) have link-distance two. Since \( x \) and \( y \) were chosen arbitrarily it follows that \( P \) is \( L \)-convex. Q.E.D.

V. A New Class of Polygons

It is interesting to consider a further generalization of Theorem 4.2 by removing condition (2) requiring that \( S^* \) be star-shaped from \( x \). Then a new class of polygons is obtained.

**Definition:** A simple polygon \( P \) is said to be \( P^*-convex \) provided that every point \( x \) in \( P \) is contained in a star-shaped subset of \( P \), from which \( P \) is weakly visible.

An \( L \)-convex polygon is \( P^*-convex \). However, the converse is no longer true as illustrated in Fig. 5. The polygon in Fig. 5 is not \( L \)-convex because the link-distance between vertices 2 and 5 is three. On the other hand the polygon is \( P^*-convex \). To see this let \( S_i \) denote the union of \( S_i \) and \( S_j \), and let \( S_2 \) denote the union of \( S_2 \) and \( S_3 \). Every point \( x \) in \( P \) must lie in either region \( S_1 \), or \( S_2 \), both regions are star-shaped from vertices 4 and 1, respectively, and \( P \) is weakly visible from each such region.

![Fig. 5 A polygon that is not L-convex but is P*-convex](image-url)

As a consequence, if a polygon is \( P^*-convex \) it must be \( L \)-convex. To see this choose any two points \( p, q \) in a polygon that is \( P^*-convex \) and let \( S^*(p) \) be the star-shaped region in \( P \) that contains \( p \) as guaranteed by the definition. Let \( q' \) be a point in \( S^*(p) \) that is visible from \( q \) as guaranteed by the definition of \( S^*(p) \). Finally, let \( k \) be a point contained in the kernel of \( S^*(p) \). Then it follows that the path \( p, k, q', q \) lies in...
\[ P \text{ and is of link-distance three. Since the choice of } p \text{ and } q \text{ was arbitrary it implies that } P \text{ is } L_3\text{-convex. On the other hand, an } L_2\text{-convex polygon is not necessarily } P^*\text{-convex, as illustrated in Fig. 6. Consider the point } p. \text{ There is no star-shaped region } S^*(p) \text{ from which } P \text{ is weakly visible. For } S^* \text{ to contain } p \text{ the kernel of } S^*(p) \text{ must lie in triangle } pxq. \text{ If this kernel lies below } [sx] \text{ then } q' \text{ is not visible from } S^*(p). \text{ On the other hand if the kernel lies above } [sx'] \text{ and close enough to } r \text{ so that } q' \text{ is visible from } S^*(p) \text{ then } r' \text{ becomes invisible from } S^*(p). \text{ Therefore following result is established.} \]

\[ \text{Fig. 6 An } L_3\text{-convex polygon that is not } P^*\text{-convex} \]

**Theorem 5.1:** \(P^*\text{-convex polygons subsume } L_2\text{-convex polygons and are a subclass of } L_1\text{-convex polygons.}

Fig. 7 illustrates the various relationships that exist between the different classes of polygons.

\[ \text{Fig. 7 The hierarchy of polygons induced by the properties of star-shapedness, } L_2\text{-convexity, } P^*\text{-convexity, and } L_1\text{-convexity} \]

VI. CONCLUSION

It would be interesting to explore the type of characterizations introduced here, based on weak-visibility from internal chords and star-shaped sets, in the case of restricted versions of convex bodies, such as shapes of constant width, as well as for convex, \(L\text{-convex and star-shaped polyhedra in three and higher dimensions.} \]

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