# An Iterative Method for the Symmetric Arrowhead Solution of Matrix Equation 

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#### Abstract

In this paper, according to the classical algorithm $L S Q R$ for solving the least-squares problem, an iterative method is proposed for least-squares solution of constrained matrix equation. By using the Kronecker product, the matrix-form $L S Q R$ is presented to obtain the like-minimum norm and minimum norm solutions in a constrained matrix set for the symmetric arrowhead matrices. Finally, numerical examples are also given to investigate the performance.


Keywords-Symmetric arrowhead matrix, iterative method, like-minimum norm, minimum norm, Algorithm $L S Q R$.

## I. INTRODUCTION

LET $R^{m \times n}$ be the set of $m \times n$ real matrices, $S A R^{n \times n}$ be the set $J_{\text {of }} n \times n$ real symmetric arrowhead matrices and $I_{n}$ be the identity matrix of order $n$. For any $A \in R^{m \times n}, A^{T}, A^{\dagger},\|A\|_{\mathrm{F}}$ and $\|A\|_{2}$ denote the transpose, Moore-Penrose generalized inverse, Frobenius norm and Euclid norm, respectively.

For $A=\left(x_{i j}\right) \in R^{m \times n}$ and $B \in R^{n \times s}, \quad A \otimes B=\left(x_{i j} B\right) \in R^{m n \times n s}$ represents the Kronecker product of $A$ and $B$.

$$
\text { For } A=\left(x_{1}, \cdots, x_{n}\right) \in R^{n \times n}, \text { define } \operatorname{vec}(A)=\left(x_{1}^{T}, x_{2}^{T}, \cdots, x_{n}^{T}\right)^{T}
$$

and $x_{\alpha: \beta, i}$ as the sub-vector consisting of the elements form $\alpha$ th component to $\beta$ th component of $x_{i}$. The inverse mapping of $\operatorname{vec}(\cdot)$ form $R^{m n}$ to $R^{m \times n}$ which is denoted by $\operatorname{mat}(\cdot)$ satisfies $\operatorname{mat}(\operatorname{vec}(A))=A$.

Definition 1. Let $A \in S A R^{n \times n}$. A is called the symmetric arrowhead matrix if it has the following form:

$$
A=\left(\begin{array}{ccccc}
x_{11} & x_{21} & x_{31} & \cdots & x_{n 1} \\
x_{21} & x_{22} & 0 & \cdots & 0 \\
x_{31} & 0 & x_{33} & \cdots & 0 \\
\vdots & \vdots & \vdots & \ddots & \vdots \\
x_{n 1} & 0 & 0 & \cdots & x_{n n}
\end{array}\right)
$$

and $\operatorname{vec}_{i}(A)$ stands for the corresponding vector

[^0]$$
\left(x_{1: n, 1}^{T}, x_{22}, x_{33}, \cdots, x_{n n}\right)^{T}
$$

Symmetric arrowhead matrix has many applications in modern control theory which can represent the parameter matrix of nonlinear control systems. Such a matrix was described as radiationless transition in the isolated molecules and oscillator attached to a Fermi liquid [1]. At present, their potential applications in electromagnetic compatibility have been more important such as the mathematical representation of interference factor.

Based on the study of [2], we consider the matrix equation

$$
\begin{equation*}
A X B+C Y D=E \tag{1}
\end{equation*}
$$

Many people have studied the matrix equation above and other constrained matrix equations, see [3], [4], [6], [8], etc. Xu, Wei, and Zhang [2] gave the solution of (1) by making use of the canonical correlation decomposition (CCD). Liao, Bai, and Lei [5] studied the least-squares solution of (1) with the least norm by combining CCD and general singular value decomposition (GSVD).

In this paper, we discuss the least-squares solution of (1) for the symmetric arrowhead matrix. When $A \in R^{m \times n}, B \in R^{n \times s}$, $C \in R^{m \times k}, D \in R^{k \times s}, E \in R^{m \times s}$ and

$$
S_{w}=\left\{[X, Y] \mid X \in S A R^{n \times n}, Y \in S A R^{k \times k}\right\}
$$

finding out $[X, Y] \in S_{w}$, such that

$$
\begin{equation*}
\min \|A X B+C Y D-E\|_{F} \tag{2}
\end{equation*}
$$

In [7], by using Moore-penrose inverse and the Kronecker product, it discussed the best approximation problem (2) and obtained a general expression of solutions.

For $H=\operatorname{diag}\left(H_{n}, H_{k}\right), P_{1}=\left(B^{T} \otimes A\right) H_{n}, P_{2}=\left(B^{T} \otimes A\right) H_{k}$, and finding $R=\left(I-P_{1} P_{1}^{+}\right) P_{2}$,

$$
\begin{aligned}
& G=R^{+}+\left(I-R^{+} R\right) Z P_{2}^{T}\left(P_{1}^{+}\right)^{T} P_{1}^{+}\left(I-P_{2} R^{+}\right), \\
& Z=\left(I+\left(I-R^{+} R\right) P_{2}^{T}\left(P_{1}^{+}\right)^{T} P_{1}^{+} P_{2}\left(I-R^{+} R\right)\right)^{-1}, \\
& K_{11}=I-P_{1}^{+} P_{1}+P_{1}^{+} P_{2} Z\left(I-R^{+} R\right) P_{2}^{T}\left(P_{1}^{+}\right)^{T}, \\
& K_{12}=-P_{1}^{+} P_{2}\left(I-R^{+} R\right) Z, K_{22}=\left(I-R^{+} R\right) Z,
\end{aligned}
$$

then the set of solutions $S_{W}$ was expressed as

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$$
S_{W}=\left\{[X, Y] \|\left[\begin{array}{c}
\operatorname{vec}(X) \\
\operatorname{vec}(Y)
\end{array}\right]=H\left[\begin{array}{c}
P_{1}^{+}-P_{1}^{+} P_{2} G \\
G
\end{array}\right] \operatorname{vec}(E)+H\left[\begin{array}{ll}
K_{11} & K_{12} \\
K_{12}^{T} & K_{13}
\end{array}\right] y\right\}
$$

where $y$ is an arbitrary vector with the proper dimension. However, the huge computation cannot be easy to realize in the large scale system which motivates us to find an operable iterative method.

A matrix pair $[X, Y]$ is referred to a minimum norm solution if it minimizes

$$
\begin{equation*}
\|X\|_{F}^{2}+\|Y\|_{F}^{2} \tag{3}
\end{equation*}
$$

and a like-minimum norm solution if it minimizes

$$
\begin{equation*}
\|\operatorname{tril}(X)\|_{F}^{2}+\|\operatorname{tril}(Y)\|_{F}^{2}, \tag{4}
\end{equation*}
$$

where $\operatorname{tril}(X)$ is denoted as lower triangular part of $X$, that is

$$
\operatorname{tril}(X)=\left(\begin{array}{cccc}
x_{11} & 0 & \cdots & 0 \\
x_{21} & x_{22} & \cdots & 0 \\
\vdots & \vdots & \ddots & \vdots \\
x_{n 1} & x_{n 2} & \cdots & x_{n n}
\end{array}\right) .
$$

## II. Preliminaries

To study the problem (2), we begin with the following lemma and the classical Algorithm $L S Q R$ presented for solving least-squares problem [3].
Lemma 1. Let $X \in S A R R^{n \times n}$, then $\operatorname{vec}(X)=\operatorname{Hvec}_{i}(X)$, where

$$
\begin{aligned}
& H_{n}=\left(\begin{array}{ccccccccccc}
e_{1} & e_{2} & e_{3} & \cdots & e_{n-1} & e_{n} & 0 & 0 & \cdots & 0 & 0 \\
0 & e_{1} & 0 & \cdots & 0 & 0 & e_{2} & 0 & \cdots & 0 & 0 \\
0 & 0 & e_{1} & \cdots & 0 & 0 & 0 & e_{3} & \cdots & 0 & 0 \\
\vdots & \vdots & \vdots & \ddots & \vdots & \vdots & \vdots & \vdots & \ddots & \vdots & \vdots \\
0 & 0 & 0 & \cdots & e_{1} & 0 & 0 & 0 & \cdots & e_{n-1} & 0 \\
0 & 0 & 0 & \cdots & 0 & e_{1} & 0 & 0 & \cdots & 0 & e_{n}
\end{array}\right), \\
& H_{n} \in R^{n^{2} \times(2 n-1)} \text { and } e_{i}=(\underbrace{0, \cdots ; 0,1, \underbrace{0}_{n-i} \cdots}_{i-1})^{T} .
\end{aligned}
$$

First, let us review the Algorithm $L S Q R$ for solving the least-squares problem:

$$
\begin{equation*}
\min _{\varphi \in R_{n}^{n}}\|M \varphi-f\|_{2} \tag{5}
\end{equation*}
$$

with given $M \in R^{m \times n}$ and vector $f \in R^{m}$, whose normal equation is

$$
\begin{equation*}
M^{T} M \varphi=M^{T} f . \tag{6}
\end{equation*}
$$

The algorithm is summarized as follows.
Algorithm LSQR
(1) Initialization

$$
\begin{aligned}
& \beta_{1} u_{1}=f, \alpha_{1} v_{1}=M^{T} u_{1}, h_{1}=v_{1}, \\
& x_{0}=0, \bar{\zeta}_{1}=\beta_{1}, \bar{\rho}_{1}=\alpha_{1} ;
\end{aligned}
$$

(2) Iteration. For $i=1,2 \ldots$
(i) bidiagonalization
(a) $\beta_{i+1} u_{i+1}=M v_{i}-\alpha_{i} u_{i}$,
(b) $\alpha_{i+1} v_{i+1}=M^{T} u_{i+1}-\beta_{i+1} v_{i}$;
(ii) construct and use Givens rotation:

$$
\begin{aligned}
& \rho_{i}=\sqrt{\bar{\rho}_{i}^{2}+\beta_{i+1}^{2}}, \\
& c_{i}=\bar{\rho}_{i} / \rho_{i}, s_{i}=\beta_{i+1} / \rho_{i}, \theta_{i+1}=s_{i} \alpha_{i+1}, \\
& \bar{\rho}_{i+1}=-c_{i} \alpha_{i+1}, \zeta_{i}=c_{i} \overline{\zeta_{i}}, \bar{\zeta}_{i+1}=s_{i} \bar{\zeta}_{i} ;
\end{aligned}
$$

(iii) update $x$ an $h$

$$
\begin{aligned}
& \varphi_{i}=\varphi_{i-1}+\left(\zeta_{i} / \rho_{i}\right) h_{i} \\
& h_{i+1}=v_{i+1}-\left(\theta_{i+1} / \rho_{i}\right) h_{i}
\end{aligned}
$$

(iv) Check convergence.

We can choose

$$
\begin{equation*}
\left|M^{T}\left(f-M \varphi_{k}\right) \|_{2}=\left|\alpha_{k+1} \bar{\zeta}_{k+1} c_{k}\right|<\tau\right. \tag{7}
\end{equation*}
$$

as convergence criteria, where $\tau>0$ is a small tolerance. Note that if (5) has a solution $\varphi \in R\left(M^{T} M\right) \in R\left(M^{T}\right)$, then $\varphi$ which is generated by Algorithm $L S Q R$ is the minimum norm solution of (5). Then we can have the symmetric arrowhead matrix solution generated by Algorithm LSQR which is the like-minimum norm and minimum norm solution of (2).

## III. The MAtrix-Form Algorithm $L S Q R$ for (2)

A. An Iterative Method for the Like-Minimum Norm Solution of (2)
In this section, we will give some results of this paper and propose iterative methods based on Algorithm LSQR.
Theorem 1. Let $X \in S A R^{n \times n}$, and $\operatorname{vec}(X)=\operatorname{Hvec}_{i}(X)$, then $H^{T} H H^{\dagger}=H^{T}$.
Proof: It follows Lemma 1,

$$
\begin{aligned}
H^{T} H & =\left(\begin{array}{ccccc}
e_{1}^{T} & 0 & \cdots & 0 & 0 \\
e_{2}^{T} & e_{1}^{T} & \cdots & 0 & 0 \\
\vdots & \vdots & & \vdots & \vdots \\
e_{n-1}^{T} & 0 & \cdots & e_{1}^{T} & 0 \\
e_{n}^{T} & 0 & \cdots & 0 & e_{1}^{T} \\
0 & e_{2}^{T} & \cdots & 0 & 0 \\
\vdots & \vdots & & \vdots & \vdots \\
0 & 0 & \cdots & e_{n-1}^{T} & 0 \\
0 & 0 & \cdots & 0 & e_{n}^{T}
\end{array}\right)\left(\begin{array}{ccccccccc}
e_{1} & e_{2} & \cdots & e_{n-1} & e_{n} & 0 & \cdots & 0 & 0 \\
0 & e_{1} & \cdots & 0 & 0 & e_{2} & \cdots & 0 & 0 \\
\vdots & \vdots & & \vdots & \vdots & \vdots & & \vdots & \vdots \\
0 & 0 & \cdots & e_{1} & 0 & 0 & \cdots & e_{n-1} & 0 \\
0 & 0 & \cdots & 0 & e_{1} & 0 & \cdots & 0 & e_{n}
\end{array}\right) \\
& =\left(\begin{array}{llllllll}
1 & & \\
& 2 I_{n-1} & \\
& & I_{n-1}
\end{array}\right) \in R^{(2 n-1) \times(2 n-1)},
\end{aligned}
$$

We can obtain matrix $H$ is full column rank, with

$$
\left(H^{T} H H^{\dagger}\right)^{T}=\left(H H^{\dagger}\right)^{T}\left(H^{T}\right)^{T}=H H^{\dagger} H=H \text {. } \square
$$

Theorem 2. Let $U \in R^{m \times s}$ and $P=A^{T} U B^{T}, Q=\frac{\left(P+P^{T}\right)}{2}$. Then for the symmetric arrowhead constrained matrix H , we have
$\left(\left(B^{T} \otimes A\right) H\right)^{T} \operatorname{vec}(U)=\left(H^{T} H\right) \operatorname{vec}_{i}\left(E_{11} Q+\left(E_{11} Q\right)^{T}+\operatorname{diag}(Q)-2 E_{11} Q E_{11}\right)$
where $E_{11}=e_{1} e_{1}^{T}$.
Proof. Notice that

$$
\begin{aligned}
& \left(\left(B^{T} \otimes A\right) H\right)^{T} \operatorname{vec}(U)=H^{T}\left(B \otimes A^{T}\right) \operatorname{vec}(U)=H^{T} \operatorname{vec}\left(A^{T} U B^{T}\right) \\
& =\left(H^{T} H\right) H^{\dagger} \operatorname{vec}\left(A^{T} U B^{T}\right)=\left(H^{T} H\right) H^{\dagger} \operatorname{vec}(P) .
\end{aligned}
$$

From Theorem 1, for any $P \in R^{n \times n}$, we have

$$
\begin{aligned}
& H^{\dagger} \operatorname{vec}(P)=\left(H^{T} H\right)^{-1} H^{T}\left(p_{11}, \cdots, p_{n 1}, p_{12}, \cdots, p_{n 2}, \cdots, p_{1 n}, \cdots, p_{n n}\right)^{T} \\
& =\left(H^{T} H\right)^{-1}\left(p_{11}, p_{12}+p_{21}, p_{13}+p_{31}, \cdots, p_{1 n}+p_{n 1}, p_{22}, p_{33}, \cdots, p_{n n}\right)^{T}
\end{aligned}
$$

that is

$$
H^{\dagger} v e c(P)=H^{\dagger} v e c\left(P^{T}\right)=H^{\dagger} v e c\left(\frac{P+P^{T}}{2}\right) .
$$

It is easy to obtain that

$$
\begin{aligned}
& H^{\dagger} \operatorname{vec}\left(\frac{P+P^{T}}{2}\right)=H^{\dagger} \operatorname{vec}\left(E_{11} Q+\left(E_{11} Q\right)^{T}+\operatorname{diag}(Q)-2 E_{11} Q E_{11}\right) \\
& =\operatorname{vec}_{i}\left(E_{11} Q+\left(E_{11} Q\right)^{T}+\operatorname{diag}(Q)-2 E_{11} Q E_{11}\right) .
\end{aligned}
$$

From all above, we can have

$$
\left(\left(B^{T} \otimes A\right) H\right)^{T} \operatorname{vec}(U)=\left(H^{T} H\right) \operatorname{vec}_{i}\left(E_{11} Q+\left(E_{11} Q\right)^{T}+\operatorname{diag}(Q)-2 E_{11} Q E_{11}\right) \text {. }
$$

Since

$$
\begin{aligned}
& \|A X B+C Y D-E\|_{F}^{2}=\left\|\left(B^{T} \otimes A, D^{T} \otimes C\right)\binom{\operatorname{vec}(X)}{\operatorname{vec}(Y)}-\operatorname{vec}(E)\right\|_{F}^{2} \\
& =\left\|\left(\left(B^{T} \otimes A\right) H_{1},\left(D^{T} \otimes C\right) H_{2}\right)\binom{\operatorname{vec}_{i}(X)}{\operatorname{vec}_{i}(Y)}-\operatorname{vec}(E)\right\|_{F}^{2},
\end{aligned}
$$

the symmetric arrowhead constrained problem is equivalent to (5) and

$$
\begin{gather*}
M=\left(\left(B^{T} \otimes A\right) H_{1},\left(D^{T} \otimes C\right) H_{2}\right) \in R^{m \times \times(2 n+2 k-2)},  \tag{8}\\
f=\operatorname{vec}(E) \in R^{m s}, \varphi=\binom{\operatorname{vec}_{i}(X)}{\operatorname{vec}_{i}(Y)} \in R^{2 n+2 k-2}, \tag{9}
\end{gather*}
$$

where $H_{1}$ and $H_{2}$ are the symmetric arrowhead constrained matrices of degree $n$ and $k$, respectively. Therefore, the normal equation of (2) is

$$
M^{T} M \varphi=M^{T} f=M^{T} \operatorname{vec}(E)
$$

Now, we will apply Algorithm $L S Q R$ to (2) and the iterative vector will be transformed into matrix so that Kronecker product and constrained matrix $H$ can be released. Then the vector $u$ and $v$ will be expressed by matrix $U$ and $V$ respectively so as to transform the matrix-vector product of $M v$ and $M^{T} u$ to the matrix-matrix form.
Let $u=v e c(U) \in R^{m s}$ with $U \in R^{m \times s}, v=\binom{v_{1}}{v_{2}} \in R^{2 n+2 k-2}$, where $v_{1}=\operatorname{vec}_{i}\left(V_{1}\right)$ and $v_{2}=\operatorname{vec}_{i}\left(V_{2}\right)$ with $V_{1} \in S A R^{n \times n}, \quad V_{2} \in S A R^{k \times k}$. Denoted by

$$
W=E_{11} Q+\left(E_{11} Q\right)^{T}+\operatorname{diag}(Q)-2 E_{11} Q E_{11},
$$

and according to Theorem 2, we have

$$
\begin{aligned}
& \operatorname{mat}\left(M^{T} \operatorname{vec}(u)\right)=\operatorname{mat}\left(H^{T} \operatorname{Hvec}_{i}\left(E_{11} Q+\left(E_{11} Q\right)^{T}+\operatorname{diag}(Q)-2 E_{11} Q E_{11}\right)\right) \\
& =\operatorname{mat}\left(\left(\begin{array}{lll}
1 & & \\
& 2 I_{n-1} & \\
& & I_{n-1}
\end{array}\right) \operatorname{vec}_{i}(W)\right)=2 W-\operatorname{diag}(W) .
\end{aligned}
$$

Then

$$
\begin{aligned}
& \operatorname{mat}\left(M^{T} u\right)=\operatorname{mat}\binom{H_{1}^{T}\left(B \otimes A^{T}\right) \operatorname{vec}(U)}{H_{2}^{T}\left(D \otimes C^{T}\right) \operatorname{vec}(U)}=\operatorname{mat}\binom{H_{1}^{T} H_{1} H_{1}^{\dagger} \operatorname{vec}\left(P_{1}\right)}{H_{2}^{T} H_{2} H_{2}^{\top} \operatorname{vec}\left(P_{2}\right)} \\
& =\operatorname{mat}\binom{H_{1}^{T} H_{1} \operatorname{vec} c_{i}\left(W_{1}\right)}{H_{2}^{T} H_{2} \operatorname{vec} c_{i}\left(W_{2}\right)}=\binom{2 W_{1}-\operatorname{diag}\left(W_{1}\right)}{2 W_{2}-\operatorname{diag}\left(W_{2}\right)} \\
& \operatorname{mat}(M v)=\operatorname{mat}\left(\left(\left(B^{T} \otimes A\right) H_{1},\left(D^{T} \otimes C\right) H_{2}\right)\binom{v_{1}}{v_{2}}\right) \\
& =\operatorname{mat}\left(\left(B^{T} \otimes A\right) H_{1} \operatorname{vec}\left(V_{1}\right)+\left(D^{T} \otimes C\right) H_{2} \operatorname{vec}_{i}\left(V_{2}\right)\right) \\
& =\operatorname{mat}\left(\left(B^{T} \otimes A\right) \operatorname{vec}\left(V_{1}\right)+\left(D^{T} \otimes C\right) \operatorname{vec}\left(V_{2}\right)\right) \\
& =A V_{1} B+C V_{2} D,
\end{aligned}
$$

where $v_{1}$ and $v_{2}$ can be formally transformed to symmetric matrices:

$$
\begin{aligned}
V_{1} & =2 W_{1}-\operatorname{diag}\left(W_{1}\right), \\
V_{2} & =2 W_{2}-\operatorname{diag}\left(W_{2}\right) .
\end{aligned}
$$

Next, we will give the algorithm for the like-minimum norm solution of (2).
Algorithm LSQR-W. 1

$$
\begin{aligned}
& \text { (1) Initialization } \\
& X_{0}=0\left(\in S A R^{n \times n}\right), Y_{0}=0\left(\in S A R^{k \times k}\right), \\
& \beta_{1}=\|E\|_{F}, U_{1}=\frac{E}{\beta_{1}}, \\
& P_{1}^{(1)}=A^{T} U_{1} B^{T}, P_{1}^{(2)}=C^{T} U_{1} D^{T} \text {, } \\
& Q_{1}^{(1)}=\frac{P_{1}^{(1)}+P_{1}^{(1)^{T}}}{2}, Q_{1}^{(2)}=\frac{P_{1}^{(2)}+P_{1}^{(2)^{T}}}{2} \text {, }
\end{aligned}
$$

$$
\begin{aligned}
& W_{1}^{(1)}=E_{11} Q_{1}^{(1)}+\left(E_{11} Q_{1}^{(1)}\right)^{T}+\operatorname{diag}\left(Q_{1}^{(1)}\right)-2 E_{11} Q_{1}^{(1)} E_{11}, \\
& W_{1}^{(2)}=E_{11} Q_{1}^{(2)}+\left(E_{11} Q_{1}^{(2)}\right)^{T}+\operatorname{diag}\left(Q_{1}^{(2)}\right)-2 E_{11} Q_{1}^{(2)} E_{11}, \\
& \bar{V}_{1}^{(1)}=2 W_{1}^{(1)}-\operatorname{diag}\left(W_{1}^{(1)}\right), \\
& \bar{V}_{1}^{(2)}=2 W_{1}^{(2)}-\operatorname{diag}\left(W_{1}^{(2)}\right), \\
& \alpha_{1}=\sqrt{\left\|\operatorname{tril}\left(\bar{V}_{1}^{(1)}\right)\right\|_{F}^{2}+\left\|\operatorname{tril}\left(\bar{V}_{1}^{(2)}\right)\right\|_{F}^{2}}, \\
& V_{1}^{(1)}=\frac{\bar{V}_{1}^{(1)}}{\alpha_{1}}, V_{1}^{(2)}=\frac{\bar{V}_{1}^{(2)}}{\alpha_{1}}, \\
& Z_{1}^{(1)}=V_{1}^{(1)}, Z_{1}^{(2)}=V_{1}^{(2)}, \\
& \bar{F}_{1}=\beta_{1}, \bar{\rho}_{1}=\alpha_{1} .
\end{aligned}
$$

(2) Iteration. For $i=1,2, \cdots$
(i) Compute $U_{i+1}$ :

$$
\begin{aligned}
& \bar{U}_{i+1}=A V_{i}^{(1)} B+C V_{i}^{(2)} D-\alpha_{i} U_{i}, \\
& \beta_{i+1}=\left\|\bar{U}_{i+1}\right\|_{F}, U_{i+1}=\frac{\bar{U}_{i+1}}{\beta_{i+1}},
\end{aligned}
$$

(ii) Compute $V_{i+1}$ :

$$
\begin{aligned}
& P_{i+1}^{(1)}=A^{T} U_{i+1} B^{T}, P_{i+1}^{(2)}=C^{T} U_{i+1} D^{T}, \\
& Q_{i+1}^{(1)}=\frac{P_{i+1}^{(1)}+P_{i+1}^{(1)}}{2}, Q_{i+1}^{(2)}=\frac{P_{i+1}^{(2)}+P_{i+1}^{(2)^{T}}}{2}, \\
& W_{i+1}^{(1)}=E_{11} Q_{i+1}^{(1)}+\left(E_{11} Q_{i+1}^{(1)}\right)^{T}+\operatorname{diag}\left(Q_{i+1}^{(1)}\right)-2 E_{11} Q_{i+1}^{(1)} E_{11}, \\
& W_{i+1}^{(2)}=E_{11} Q_{1}^{(2)}+\left(E_{11} Q_{i+1}^{(2)}\right)^{T}+\operatorname{diag}\left(Q_{i+1}^{(2)}\right)-2 E_{11} Q_{i+1}^{(2)} E_{11}, \\
& \bar{V}_{i+1}^{(1)}=2 W_{i+1}^{(1)}-\operatorname{diag}\left(W_{i+1}^{(1)}\right)-\beta_{i+1} V_{i}^{(1)}, \\
& \bar{V}_{i+1}^{(2)}=2 W_{i+1}^{(2)}-\operatorname{diag}\left(W_{i+1}^{(2)}\right)-\beta_{i+1} V_{i}^{(2)}, \\
& \alpha_{i+1}=\sqrt{\left\|\operatorname{tril}\left(\bar{V}_{i+1}^{(1)}\right)\right\|_{F}^{2}+\left\|\operatorname{tril}\left(\bar{V}_{i+1}^{(2)}\right)\right\|_{F}^{2}}, \\
& V_{i+1}^{(1)}=\frac{\bar{V}_{i+1}^{(1)}}{\alpha_{i+1}}, V_{i+1}^{(2)}=\frac{\bar{V}_{i+1}^{(2)}}{\alpha_{i+1}},
\end{aligned}
$$

(iii) Compute Givens rotation:

$$
\begin{aligned}
& \rho_{i}=\sqrt{\bar{\rho}_{i}^{2}+\beta_{i+1}^{2}}, \\
& c_{i}=\frac{\bar{\rho}_{i}}{\rho_{i}}, s_{i}=\frac{\beta_{i+1}}{\rho_{i}}, \theta_{i+1}=s_{i} \alpha_{i+1}, \\
& \bar{\rho}_{i+1}=-c_{i} \alpha_{i+1}, \xi_{i}=c_{i} \bar{\xi}_{i}, \bar{\xi}_{i+1}=s_{i} \bar{\xi}_{i},
\end{aligned}
$$

(iv) Update $X_{i}, Y_{i}$ and $Z_{i}$ :

$$
\begin{aligned}
& X_{i}=X_{i-1}+\left(\xi_{i} / \rho_{i}\right) Z_{i}^{(1)}, \\
& Y_{i}=Y_{i-1}+\left(\xi_{i} / \rho_{i}\right) Z_{i}^{(2)}, \\
& Z_{i+1}^{(1)}=V_{i+1}^{(1)}-\left(\theta_{i} / \rho_{i}\right) Z_{i}^{(1)}, \\
& Z_{i+1}^{(2)}=V_{i+1}^{(2)}-\left(\theta_{i} / \rho_{i}\right) Z_{i}^{(2)} .
\end{aligned}
$$

(3) Check convergence.

Algorithm $L S Q R-W .1$ can compute the like-minimum norm solution $\varphi=\binom{\operatorname{vec}_{i}(X)}{\operatorname{vec}_{i}(Y)}$ of (5). So we have the following result.
Theorem 3. The symmetric arrowhead solution generated by Algorithm LSQR-W. 1 is the like-minimum norm solution of (2).
B. An Iterative Method for the Minimum Norm Solution of (2)

In Section $A$, the like-minimum solution generated by Algorithm $L S Q R-W .1$ is not the best approximation result and non-unique. In this section, we will give an iterative method for the minimum norm solution of (2).

For $X \in R^{n \times n}$, define $v \overline{e e}_{i}(X)=\operatorname{Svec}_{i}(X)$, and add the weight values to elements $x_{i j}$, that is

with $S \in R^{(2 n-1) \times(2 n-1)}$. Obviously, there is one to one linear mapping from $v \bar{e} c_{i}(X)$ to $\operatorname{vec}(X)$. Let $\bar{H}$ denote the minimum norm constrained matrix with

$$
\operatorname{vec}(X)=\bar{H} v \bar{e} c_{i}(X)
$$

It readily follows form Theorem 1 that
Theorem 4. Suppose $\bar{H}$ is the symmetric arrowhead constrained matrix, then

$$
\bar{H}=H S^{-1}
$$

and

$$
\bar{H}^{T} H=2 I_{2 n-1}
$$

Since

$$
\|A X B+C Y D-E\|_{F}^{2}=\left\|\left(\left(B^{T} \otimes A\right) \bar{H}_{1},\left(D^{T} \otimes C\right) \bar{H}_{2}\right)\binom{\operatorname{vec}_{i}(X)}{\operatorname{vec}_{i}(Y)}-v e c(E)\right\|_{F}^{2},
$$

problem (2) is equivalent to (5) and

$$
\begin{gathered}
M=\left(\left(B^{T} \otimes A\right) \bar{H}_{1},\left(D^{T} \otimes C\right) \bar{H}_{2}\right) \in R^{m s \times(2 n+2 k-2)}, \\
f=v e c(E) \in R^{m s}, x=\binom{v \overline{e_{c}}(X)}{v \overline{e_{i}}(X)} \in R^{2 n+2 k-2},
\end{gathered}
$$

where $\bar{H}_{1}$ and $\bar{H}_{2}$ are the new symmetric arrowhead constrained matrices of degree $n$ and $k$, respectively.
For any $v=\binom{v_{1}}{v_{2}} \in R^{2 n+2 k-2}, v_{1}=\operatorname{vec}_{i}\left(V_{1}\right)$ and $v_{2}=\operatorname{vec}_{i}\left(V_{2}\right)$ where $V_{1} \in S A R^{n \times n}, V_{2} \in S A R^{k \times k}$, we have

$$
\tilde{V}=V+(\sqrt{2}-1) \operatorname{diag}(V)
$$

and

$$
\begin{aligned}
& \operatorname{mat}(M v)=\operatorname{mat}\left(\left(B^{T} \otimes A\right) \bar{H}_{1} v e c_{i}\left(V_{1}\right)+\left(D^{T} \otimes C\right) \bar{H}_{2} \operatorname{vec}_{i}\left(V_{2}\right)\right) \\
& =\operatorname{mat}\left(\left(B^{T} \otimes A\right) H_{1} S_{1}^{-1} v \overline{v e}_{i}\left(V_{1}\right)+\left(D^{T} \otimes C\right) H_{2} S_{2}^{-1} \overline{v e}_{c_{i}}\left(V_{2}\right)\right) \\
& =\operatorname{mat}\left(\left(B^{T} \otimes A\right) H_{1} v e c_{i}\left(\tilde{V}_{1}\right)+\left(D^{T} \otimes C\right) H_{2} v e c_{i}\left(\tilde{V}_{2}\right)\right) \\
& =\operatorname{mat}\left(\left(B^{T} \otimes A\right) \operatorname{vec}\left(\tilde{V}_{1}\right)+\left(D^{T} \otimes C\right) \operatorname{vec}\left(\tilde{V}_{2}\right)\right) \\
& =A \tilde{V}_{1} B+C \tilde{V}_{2} D,
\end{aligned}
$$

For any $u=\operatorname{vec}(U) \in R^{n s}$ with $U \in R^{m \times s}$, denoted by

$$
W=E_{11} Q+\left(E_{11} Q\right)^{T}+\operatorname{diag}(Q)-2 E_{11} Q E_{11},
$$

according to Theorem 2, we have

$$
\begin{aligned}
& \operatorname{mat}\left(M^{T} u\right)=\operatorname{mat}\binom{\bar{H}_{1}^{T}\left(B \otimes A^{T}\right) \operatorname{vec}(U)}{\bar{H}_{2}^{T}\left(D \otimes C^{T}\right) \operatorname{vec}(U)}=\operatorname{mat}\binom{\bar{H}_{1}^{T} \bar{H}_{1} \bar{H}_{1}^{\dagger} \operatorname{vec}\left(P_{1}\right)}{\bar{H}_{2}^{T} \bar{H}_{2} \bar{H}_{2}^{\dagger} \operatorname{vec}\left(P_{2}\right)} \\
& =\operatorname{mat}\binom{2 I_{n} S H_{1}^{\dagger} \operatorname{vec}(W)}{2 I_{k} S S_{2}^{\dagger} \operatorname{vec}(W)} \operatorname{mat}\binom{2 \operatorname{Svec}_{i_{2}}\left(W_{1}\right)}{2 \operatorname{Svec}_{i}\left(W_{2}\right)} \\
& =\binom{2 W_{1}-(2-\sqrt{2}) \operatorname{diag}\left(W_{1}\right)}{2 W_{2}-(2-\sqrt{2}) \operatorname{diag}\left(W_{2}\right)},
\end{aligned}
$$

where $v_{1}$ and $v_{2}$ can be formally transformed to symmetric matrices

$$
\begin{gathered}
V_{1}=2 W_{1}-(2-\sqrt{2}) \operatorname{diag}\left(W_{1}\right) \\
V_{2}=2 W_{2}-(2-\sqrt{2}) \operatorname{diag}\left(W_{2}\right) .
\end{gathered}
$$

Next, we give the algorithm for the minimum norm solution of (2).

## Algorithm LSQR-W. 2

(1) Initialization

$$
\begin{aligned}
& X_{0}=0\left(\in S A R^{n \times n}\right), Y_{0}=0\left(\in S A R^{1 \times k}\right), \\
& \beta_{1}=\|E\|_{F}, U_{1}=\frac{E}{\beta_{1}}, \\
& P_{1}^{(1)}=A^{T} U_{1} B^{T}, P_{1}^{(2)}=C^{T} U_{1} D^{T}, \\
& Q_{1}^{(1)}=\frac{P_{1}^{(1)}+P_{1}^{(1)^{T}}}{2}, Q_{1}^{(2)}=\frac{P_{1}^{(2)}+P_{1}^{(2)^{T}}}{2}, \\
& W_{1}^{(1)}=E_{11} Q_{1}^{(1)}+\left(E_{11} Q_{1}^{(1)}\right)^{T}+\operatorname{diag}\left(Q_{1}^{(1)}\right)-2 E_{11} Q_{1}^{(1)} E_{11}, \\
& W_{1}^{(2)}=E_{11} Q_{1}^{(2)}+\left(E_{11} Q_{1}^{(2)}\right)^{T}+\operatorname{diag}\left(Q_{1}^{(2)}\right)-2 E_{11} Q_{1}^{(2)} E_{11}, \\
& \bar{V}_{1}^{(1)}=2 W_{1}^{(1)}-(2-\sqrt{2}) \operatorname{diag}\left(W_{1}^{(1)}\right), \\
& \bar{V}_{1}^{(2)}=2 W_{1}^{(2)}-(2-\sqrt{2}) \operatorname{diag}\left(W_{1}^{(2)}\right), \\
& \alpha_{1}=\sqrt{\left\|\operatorname{tril}\left(\bar{V}_{1}^{(1)}\right)\right\|_{F}^{2}+\left\|\operatorname{tril}\left(\bar{V}_{1}^{(2)}\right)\right\|_{F}^{2}}, \\
& V_{1}^{(1)}=\frac{\bar{V}_{1}^{(1)}}{\alpha_{1}}, V_{1}^{(2)}=\frac{\bar{V}_{1}^{(2)}}{\alpha_{1}}, \\
& Z_{1}^{(1)}=V_{1}^{(1)}, Z_{1}^{(2)}=V_{1}^{(2)}, \\
& \bar{E}_{1}=\beta_{1}, \bar{\rho}_{1}=\alpha_{1} ;
\end{aligned}
$$

(2) Iteration. For $i=1,2, \cdots$
(i) Compute $U_{i+1}$ :

$$
\begin{aligned}
& \tilde{V}_{i}^{(1)}=V_{i}^{(1)}+(\sqrt{2}-1) \operatorname{diag}\left(V_{i}^{(1)}\right), \\
& \tilde{V}_{i}^{(2)}=V_{i}^{(2)}+(\sqrt{2}-1) \operatorname{diag}\left(V_{i}^{(2)}\right), \\
& \bar{U}_{i+1}=A \tilde{V}_{i}^{(1)} B+C \tilde{V}_{i}^{(2)} D-\alpha_{i} U_{i}, \\
& \beta_{i+1}=\left\|\bar{U}_{i+1}\right\|_{F}, U_{i+1}=\bar{U}_{i+1} / \beta_{i+1},
\end{aligned}
$$

(ii) Compute $V_{i+1}$ :

$$
\begin{aligned}
& P_{i+1}^{(1)}=A^{T} U_{i+1} B^{T}, P_{i+1}^{(2)}=C^{T} U_{i+1} D^{T}, \\
& Q_{i+1}^{(1)}=\frac{P_{i+1}^{(1)}+P_{i+1}^{(1)^{T}}}{2}, Q_{i+1}^{(2)}=\frac{P_{i+1}^{(2)}+P_{i+1}^{(2)^{T}}}{2}, \\
& W_{i+1}^{(1)}=E_{11} Q_{i+1}^{(1)}+\left(E_{11} Q_{i+1}^{(1)}\right)^{T}+\operatorname{diag}\left(Q_{i+1}^{(1)}\right)-2 E_{11} Q_{i+1}^{(1)} E_{11}, \\
& W_{i+1}^{(2)}=E_{11} Q_{1}^{(2)}+\left(E_{11} Q_{i+1}^{(2)}\right)^{T}+\operatorname{diag}\left(Q_{i+1}^{(2)}\right)-2 E_{11} Q_{i+1}^{(2)} E_{11}, \\
& \bar{V}_{i+1}^{(1)}=2 W_{i+1}^{(1)}-(2-\sqrt{2}) \operatorname{diag}\left(W_{i+1}^{(1)}\right)-\beta_{i+1} V_{i}^{(1)}, \\
& \bar{V}_{i+1}^{(2)}=2 W_{i+1}^{(2)}-(2-\sqrt{2}) \operatorname{diag}\left(W_{i+1}^{(2)}\right)-\beta_{i+1} V_{i}^{(2)}, \\
& \alpha_{i+1}=\sqrt{\left\|\operatorname{tril}\left(\bar{V}_{i+1}^{(1)}\right)\right\|_{F}^{2}+\left\|\operatorname{tril}\left(\bar{V}_{i+1}^{(2)}\right)\right\|_{F}^{2}}, \\
& V_{i+1}^{(1)}=\frac{\bar{V}_{i+1}^{(1)}}{\alpha_{i+1}}, V_{i+1}^{(2)}=\frac{\bar{V}_{i+1}^{(2)}}{\alpha_{i+1}},
\end{aligned}
$$

(iii) Compute Givens rotation:

$$
\begin{aligned}
& \rho_{i}=\sqrt{\bar{\rho}_{i}^{2}+\beta_{i+1}^{2}}, \\
& c_{i}=\frac{\bar{\rho}_{i}}{\rho_{i}}, s_{i}=\frac{\beta_{i+1}}{\rho_{i}}, \theta_{i+1}=s_{i} \alpha_{i+1}, \\
& \bar{\rho}_{i+1}=-c_{i} \alpha_{i+1}, \xi_{i}=c_{i} \bar{\xi}_{i}, \bar{\xi}_{i+1}=s_{i} \bar{\xi}_{i},
\end{aligned}
$$

(iv) Update $X_{i}, Y_{i}$ and $Z_{i}$ :

$$
\begin{aligned}
& x_{i}=x_{i-1}+\left(\xi_{i} / \rho_{i}\right) Z_{i}^{(1)}, \\
& y_{i}=y_{i-1}+\left(\xi_{i} / \rho_{i}\right) Z_{i}^{(2)}, \\
& X_{i}=x_{i}+(\sqrt{2}-1) \operatorname{diag}\left(x_{i}\right), \\
& Y_{i}=y_{i}+(\sqrt{2}-1) \operatorname{diag}\left(y_{i}\right) . \\
& Z_{i+1}^{(1)}=V_{i+1}^{(1)}-\left(\theta_{i} / \rho_{i}\right) Z_{i}^{(1)}, \\
& Z_{i+1}^{(2)}=V_{i+1}^{(2)}-\left(\theta_{i} / \rho_{i}\right) Z_{i}^{(2)} .
\end{aligned}
$$

(3) Check convergence.

From above translation process and the algorithm, we can obtain the approximate solution $\left[X_{i}, Y_{i}\right]$ with

$$
\operatorname{vec}_{i}\left(x_{i}\right)=v \bar{e} c_{i}\left(X_{i}\right), v e c_{i}\left(y_{i}\right)=v \bar{e} c_{i}\left(Y_{i}\right) .
$$

Thus, we have

$$
\begin{aligned}
& X_{i}=x_{i}+(\sqrt{2}-1) \operatorname{diag}\left(x_{i}\right), \\
& Y_{i}=y_{i}+(\sqrt{2}-1) \operatorname{diag}\left(y_{i}\right) .
\end{aligned}
$$

Because of

$$
\|X\|_{F}^{2}+\|Y\|_{F}^{2}=2\left(\left\|v \overline{e 匕}_{i}(X)\right\|_{2}^{2}+\left\|v \overline{e c}_{i}(Y)\right\|_{2}^{2}\right),
$$

The Algorithm LSQR-W. 2 can compute the minimum norm solution of (5). So we have the following result.

Theorem 5. The symmetric arrowhead solution generated by Algorithm LSQR-W. 2 is the minimum norm solution of (2).

## IV. Numerical Examples

In this section, we reported three numerical examples to illustrate the efficiency of the algorithms we proposed. First, we present the Example 1 in [7] with the results generated by Algorithm LSQR-W. 1 and Algorithm LSQR-W.2.
Example 1. Considering the following matrix equation $A X B+C Y D=E$ with $A=B=C=D=I_{4}$ and

$$
E=\left[\begin{array}{llll}
1 & 1 & 1 & 1 \\
1 & 1 & 0 & 0 \\
1 & 0 & 1 & 0 \\
1 & 0 & 0 & 1
\end{array}\right]
$$

i.e.,

$$
I_{4} X I_{4}+I_{4} Y I_{4}=E
$$

where $I_{4}$ is the unit matrix of order 4 . We can also obtain the minimum solution by Algorithm LSQR-W. 1 and Algorithm $L S Q R-W .2$ which should be

$$
X=\left[\begin{array}{cccc}
0.5 & 0.5 & 0.5 & 0.5 \\
0.5 & 0.5 & 0 & 0 \\
0.5 & 0 & 0.5 & 0 \\
0.5 & 0 & 0 & 0.5
\end{array}\right], Y=\left[\begin{array}{cccc}
0.5 & 0.5 & 0.5 & 0.5 \\
0.5 & 0.5 & 0 & 0 \\
0.5 & 0 & 0.5 & 0 \\
0.5 & 0 & 0 & 0.5
\end{array}\right]
$$

The exact result could be computed by our methods in no more than two iterations. Fig. 1 plots the relation between error

$$
\gamma_{k}=\log 10\left(\left\|A X_{k} B+C Y_{k} D-E\right\|_{F}\right)
$$

and iterative number $K$. Next we present two other matrix equations which show the methods numerically reliable in various circumstances.

## Example 2. Given

$$
\begin{gathered}
A=\left[\begin{array}{cc}
\operatorname{hilb}(5) & \operatorname{zeros}(5,3) \\
\operatorname{eye}(5) & \operatorname{ones}(5,3)
\end{array}\right], B=\left[\begin{array}{cc}
\operatorname{ones}(3,7) & \operatorname{zeros}(3,5) \\
\operatorname{zeros}(5,7) & \operatorname{pascal}(5)
\end{array}\right], \\
C=\left[\begin{array}{c}
\operatorname{magic}(6) \\
\operatorname{ones}(4,6)
\end{array}\right], D=\left[\begin{array}{cc}
\operatorname{hankel}(1: 4) & \operatorname{zeros}(4,8) \\
\operatorname{zeros}(2,4) & \operatorname{ones}(2,8)
\end{array}\right], \\
X=\operatorname{ones}(8,8), Y=\operatorname{ones}(6,6) \text { and } E=A X B+C Y D .
\end{gathered}
$$

Notice that problem (1) is consistent and has a symmetric arrowhead solution. For $M$ and $\varphi$ defined by (8) and (9), we can choose residual error

$$
\left\|A X_{k} B+C Y_{k} D-E\right\|_{F}=\left\|M x_{k}-f\right\|_{2}=\bar{\xi}_{k+1} .
$$

From (3) and (4), we can compute by Algorithm LSQR-W.I and

$$
\begin{aligned}
& \|X\|_{F}^{2}+\|Y\|_{F}^{2}=38,\|\operatorname{tril}(X)\|_{F}^{2}+\|\operatorname{tril}(Y)\|_{F}^{2}=26, \\
& \left\|X_{83}\right\|_{F}^{2}+\left\|Y_{83}\right\|_{F}^{2}=38.9580, \| \operatorname{rril}\left(X_{83}\left\|_{F}^{2}+\right\| \operatorname{rril}\left(Y_{83}\right) \|_{F}^{2}=25.5309 .\right.
\end{aligned}
$$

Thus, we can obtain the like-minimum norm solution of (2) and Fig. 2 plots the relation between error

$$
\varepsilon_{k}=\log 10\left(\left\|A X_{k} B+C Y_{k} D-E\right\|_{F}\right)
$$

and iterative number $K$ which shows the favorable efficiency of Algorithm LSQR-W.1. Next, we also compute the result by Algorithm LSQR-W. 2 and obtain that

$$
\left\|X_{83}\right\|_{F}^{2}+\left\|Y_{83}\right\|_{F}^{2}=38,\left\|\operatorname{tril}\left(X_{83}\right)\right\|_{F}^{2}+\left\|\operatorname{tril}\left(Y_{83}\right)\right\|_{F}^{2}=26 .
$$

Thus, we obtain the minimum norm solution of (2) and Fig. 3 plots the relation between error

$$
\mu_{k}=\log 10\left(\left\|A X_{k} B+C Y_{k} D-E\right\|_{F}\right)
$$

and iterative number $K$ of the Algorithm LSQR-W.2.
Example 3. Given

$$
\begin{gathered}
A=\left[\begin{array}{cc}
\operatorname{hilb}(5) & \operatorname{zeros}(5,3) \\
\operatorname{eye}(5) & \operatorname{ones}(5,3)
\end{array}\right], B=\left[\begin{array}{cc}
\operatorname{ones}(3,7) & \operatorname{zeros}(3,5) \\
\operatorname{zeros}(5,7) & \operatorname{pascal}(5)
\end{array}\right], \\
C=\left[\begin{array}{c}
\operatorname{magic}(6) \\
\operatorname{ones}(4,6)
\end{array}\right], D=\left[\begin{array}{cc}
\operatorname{hankel}(1: 4) & \operatorname{zeros}(4,8) \\
\operatorname{zeros}(2,4) & \operatorname{ones}(2,8)
\end{array}\right], \\
E=[\operatorname{toeplitz}(1: 10) \\
\text { ones }(10,2)] .
\end{gathered}
$$

Notice that the problem (1) is not consistent. For $M$ and $\varphi$ defined by (8) and (9), we choose residual error

$$
\left\|M^{T} M x_{k}-M^{T} f\right\|_{2}=\left|\alpha_{k+1} \bar{\xi}_{k+1} c_{k}\right| .
$$

Then Figs. 4 and 5 plot the relation between error

$$
\eta_{k}, \delta_{k}=\log 10\left(\left\|M^{T}\left(M x_{k}-f\right)\right\|_{2}\right)
$$

and iterative number $K$ by Algorithm $L S Q R-W .1$ and Algorithm $L S Q R-W .2$, respectively.


Fig. 1 The relation between error $\gamma_{k}$ and iterative number $K$


Fig. 2 The relation between error $\varepsilon_{k}$ and iterative number $K$


Fig. 3 The relation between error $\mu_{k}$ and iterative number $K$


Fig. 4 The relation between error $\eta_{k}$ and iterative number $K$


Fig. 5 The relation between error $\delta_{k}$ and iterative number $K$

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