

Density Clustering Based on Radius of Data (DCBRD)

A.M. Fahim, A. M. Salem, F. A. Torkey, and M. A. Ramadan

Abstract—Clustering algorithms are attractive for the task of class identification in spatial databases. However, the application to large spatial databases rises the following requirements for clustering algorithms: minimal requirements of domain knowledge to determine the input parameters, discovery of clusters with arbitrary shape and good efficiency on large databases. The well-known clustering algorithms offer no solution to the combination of these requirements. In this paper, a density based clustering algorithm (DCBRD) is presented, relying on a knowledge acquired from the data by dividing the data space into overlapped regions. The proposed algorithm discovers arbitrary shaped clusters, requires no input parameters and uses the same definitions of DBSCAN algorithm. We performed an experimental evaluation of the effectiveness and efficiency of it, and compared this results with that of DBSCAN. The results of our experiments demonstrate that the proposed algorithm is significantly efficient in discovering clusters of arbitrary shape and size.

Keywords—Clustering Algorithms, Arbitrary Shape of clusters, cluster Analysis.

I. INTRODUCTION

LARGE datasets have been collected or produced in many application domains, such as bioinformatics, physics, geology, and marketing, and some have reached the level of terabytes. Since the knowledge hidden in this data is usually of great strategic importance [10]. One of the primary data analysis tasks is the cluster analysis, which help user to uncover the knowledge hidden in the collected data. Clustering is one of the most important tasks in data mining and knowledge discovery [5]. Clustering groups database data into meaningful subclasses in such a way that minimizes the intra-differences and maximizes the inter-differences of these subclasses [15]. Up to now, many clustering algorithms have been proposed, each of these algorithms have drawbacks and advantages. A clustering algorithm is considered to be good if it satisfies the following requirements, (1) minimal

requirements of domain knowledge to determine the values of its input parameters, which is very important problem especially for large data sets. (2) Discovery of arbitrary shaped clusters. (3) good efficiency on large data sets, data set may contains large number of objects or the object described by large number of features or data is large in both previous dimension. The well-known clustering algorithms offer no solution to the combination of these requirements.

In this paper we propose a clustering algorithm based on knowledge acquired from the data set, and apply the main idea of density based clustering algorithms like DBSCAN. The proposed algorithm will be called density clustering based on radius of data (DCBRD). The DCBRD algorithm requires no input parameters, discovers arbitrary size and shaped clusters, is efficient even for large data sets especially data with large dimension. The paper is organized as follows. In section 2, some previous clustering algorithms are listed. In section 3, simplified review about DBSCAN algorithm is presented. In section 4, the overlapped circular regions are presented, and we show how these circles are created and how the algorithm uses them. Section 5 demonstrates the proposed algorithm. We describe the experimental results in section 6 and conclude with section 7.

II. RELATED WORK

There are many clustering algorithms proposed, these algorithms may be classified into partitioning, hierarchical, density and grid based methods [8]. The first two types are the most common. Partitioning algorithms are k-means [12] and k-medoid [11], [13]. Hierarchical algorithms create a hierarchical decomposition of a database D . The basic hierarchical clustering algorithm works as in [6]. Some hierarchical algorithms are single-link [14], complete-link [2], [9], average-link method [16], BIRCH [17] and CURE [7].

Density-Based Clustering algorithms group objects according to specific density objective functions. The most popular one is probably DBSCAN [3]. In this paper new index structure is proposed. This structure also used to obtain an optimal value for the radius of the neighborhood. The proposed algorithm is based on this structure and the main idea of DBSCAN algorithm.

III. DBSCAN ALGORITHM

The key idea of density-based clustering is that for each object of a cluster the neighborhood of a given radius (Eps) has to contain at least a minimum number of objects ($MinPts$), i.e. the cardinality of the neighborhood has to exceed some threshold. We will first give a short introduction to DBSCAN. For details see [3].

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Ahmed Mahmoud Fahim is PhD student in the faculty of Education, Suez Canal University, Egypt (e-mail: ahmedfahim@yahoo.com).

Abdel-Badeeh Mohamed Salem is with Ain shams University, Faculty of Computers and Information, Computer Science Department, Egypt (e-mail: absalem@asunet.shams.edu.eg).

Fawzy Aly Torkey is the Dean of Faculty of Computers and Information, Menoufia University, Egypt (e-mail: fatorkey@yahoo.com).

Mohamed A. Ramadan is with Faculty of Science, Mathematic Department, Menoufia University, Egypt (e-mail: mramadan@mail.eun.eg).

Definition 1: (directly density-reachable) An object p is *directly density-reachable* from an object q wrt. Eps and $MinPts$ in the set of objects D if :

1) $p \in N_{Eps}(q)$, ($N_{Eps}(q)$ is the subset of D contained in the Eps -neighborhood of q).

2) $Card(N_{Eps}(q)) \geq MinPts$.

Definition 2: (density-reachable) An object p is *density-reachable* from an object q wrt. Eps and $MinPts$ in the set of objects D , denoted as $p > Dq$, if there is a chain of objects p_1, \dots, p_m $p_1 = q, p_m = p$ such that $p_i \in D$ and p_{i+1} is directly density-reachable from p_i wrt. Eps and $MinPts$.

Definition 3: (density-connected) An object p is *density-connected* to an object q wrt. Eps and $MinPts$ in the set of objects D if there is an object $o \in D$ such that both p and q are density-reachable from o wrt. Eps and $MinPts$ in D .

Fig. 1 illustrates the definitions on a sample database of objects from a 2-dimensional vector space.

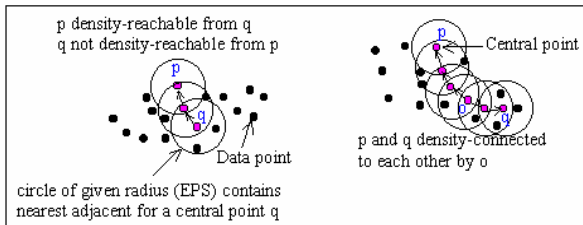


Fig. 1 Density-reachability and density-connectivity

Definition 4: (cluster) Let D be a set of objects. A *cluster* C wrt. Eps and $MinPts$ in D is a non-empty subset of D satisfying the following conditions:

1) **Maximality:** $\forall p, q \in D$: if $p \in C$ and $q > D p$ wrt. Eps and $MinPts$, then also $q \in C$.

2) **Connectivity:** $\forall p, q \in C$: p is density-connected to q wrt. Eps and $MinPts$ in D .

Definition 5: (noise) Let C_1, \dots, C_k be the clusters wrt. Eps and $MinPts$ in D . Then, we define the *noise* as the set of objects in the database D not belonging to any cluster C_i , i.e. $noise = \{p \in D \mid \forall i: p \notin C_i\}$.

To find a cluster, DBSCAN starts with an arbitrary object p in D and retrieves all objects of D density-reachable from p with respect to Eps and $MinPts$. If p is a border object, no objects are density-reachable from p and p is assigned to noise temporarily. Then DBSCAN handles the next object in database D . Retrieval of density-reachable objects is performed by successive region queries. A region query returns all objects intersecting a specified query region efficiently by R*-trees. Before clustering the database, R*-tree should be built in advance. However, there are some DBSCAN algorithm problems limiting its applications. The DBSCAN algorithm is sketched in Fig. 2.

```

Algorithm DBSCAN ( $D, Eps, MinPts$ )
// All objects in  $D$  are unclassified.
begin
  FORALL objects  $o$  in  $D$  DO:
    IF  $o$  is unclassified
      call function expand_cluster to construct a cluster wrt.  $Eps$ 
      and  $MinPts$  containing  $o$ .
    end
  end

  FUNCTION expand_cluster ( $o, D, Eps, MinPts$ ):
  begin
    retrieve the  $Eps$ -neighborhood  $N_{Eps}(o)$  of  $o$ ;
    IF  $|N_{Eps}(o)| \geq MinPts$ , // i.e.  $o$  is not a core object
      mark  $o$  as noise and RETURN;
    ELSE // i.e.  $o$  is a core object
      select a new cluster-id and mark all objects in  $N_{Eps}(o)$  with
      this current cluster-id;
      push all objects from  $N_{Eps}(o)$  onto the stack seeds;
      WHILE NOT seeds.empty() DO
        currentObject := seeds.top();
        retrieve the  $Eps$ -neighborhood  $N_{Eps}(currentObject)$  of
        currentObject;
        IF  $|N_{Eps}(currentObject)| \geq MinPts$ 
          select all objects in  $N_{Eps}(currentObject)$  not yet classified or
          are marked as noise;
          push the unclassified objects onto seeds
          and mark all of these objects with current cluster-id;
        seeds.pop();
      END
    RETURN
  end

```

Fig. 2 DBSCAN Algorithm

The most fundamental open problem is: DBSCAN requires the user to specify a global threshold Eps ($Minpts$ is often fixed to 4 to reduce the computational amount). In order to determine Eps , DBSCAN has to calculate the distance between an object and its k th ($k=4$) nearest neighbor for all objects. In addition, DBSCAN is based on R*-tree, and calculates the k -dist value on the entire database. The two procedures are the most time-consuming phases in the whole clustering process, but their computational loads are not included in time consumption as in $O(n \log n)$, so the actual time consumption of DBSCAN may be larger than that of $O(n \log n)$ [15]. Clustering procedure is very expensive so that it is computationally prohibitive for large databases. Eps and $Minpts$ determine a density threshold, thus DBSCAN becomes a typical density-based clustering method. Furthermore, the $Minpts$ usually is fixed to 4, thus the density threshold is perfectly determined by Eps .

IV. DATA SPACE PARTITIONING INTO OVERLAPPED CIRCLES

The key idea of the proposed algorithm is that one can greatly reduce the number of distance computations required for clustering by partitioning the data into overlapping subsets, and then only measuring distances among pairs of data points that belong to a common subset. The overlapped circle technique thus uses two different sources of information to cluster items: a cheap and approximate similarity measure (the radius of the overlapped circles that cover all data space) and a more expensive and accurate similarity measure (the optimal value of Eps).

The proposed algorithm divides the clustering process into two stages. In the first stage, the algorithm use the cheap distance measure in order to creates some number of overlapping circles (or hyper sphere), any circle is simply a subset of the elements (i.e. data points or items) that, according to the approximate similarity measure, are within some distance threshold from a central point. Significantly, a data point may appear under more than one circle, and every data point must appear in at least one circle. The circles with

solid outlines in Fig. 3 show an example of overlapping circles that cover a data set. Note that every solid circle (contains nearest adjacent for a central point) is a subset of larger dashed circle (contains nearest far adjacent for a central point). Dashed circles are used to ensure that no cluster split into more than one cluster. If you look at cluster 1, this cluster split into two clusters without dashed circle because there is no data point in the intersection of the two inner solid circles, using the dashed circle when it is required, the algorithm detects the actual clusters contained in data set.

In the second stage, the proposed algorithm executes DBSCAN clustering algorithm, the value of Eps is obtained from the overlapped circular regions, thus the proposed algorithm does not require any input parameter.

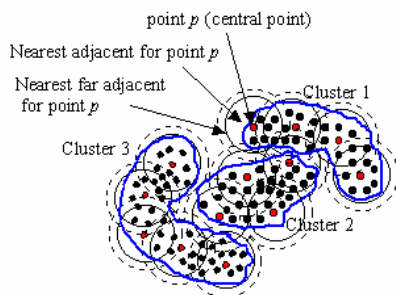


Fig. 3 An example of three clusters and the overlapped circles cover them

A. Creation of Overlapped Circles

The proposed algorithm computes the radius (Rad) of circles that cover all data points as we will see in the next section. Fig. 4 illustrates the creation of overlapped circles.

```

Function Creat_circles( $D, Rad$ )
// creat all required overlapped circles that cover all data points
// all points in  $D$  are uncovered
begin
1.  $circle\_id = 0$ 
2.  $circle\_center[circle\_id] =$  the first point in  $D$ 
3.  $next\_center = false$ 
4. for all point  $p$  in  $D$ 
5. if ( $distance(p, circle\_center[circle\_id]) \leq Rad$ )
6. { push point  $p$  to linked list1 of the  $circle\_id$ 
7.   if  $p.distance > distance(p, circle\_center[circle\_id])$ 
8.      $p.distance = distance(p, circle\_center[circle\_id])$ 
9.      $p.circle = circle\_id$ 
10. }
11. else if ( $distance(p, circle\_center[circle\_id]) \leq Rad * 1.5$ )
12.   push point  $p$  to linked list2 of the  $circle\_id$ 
13. else if ( $distance(p, circle\_center[circle\_id]) \leq Rad * 2$ )
14.   and ( $p$  is uncovered)
15.   {  $circle\_center[circle\_id + 1] = p$ 
16.      $next\_center = true$ 
17.   }
18. if ( $next\_center == true$ )
19.    $circle\_id++$ , goto step 3
20. endfor
21. for all points  $q$  in  $D$ 
22. if  $q$  is uncovered point
23.    $circle\_center[circle\_id++] = q$ , goto step 3
end function

```

Fig. 4 Creation of circles that cover all data points

The function takes the first point as the center of the first circle (step 2 of Fig. 4), and assign all points whose distances

from this center are less than or equal to the value of Rad to the list1 (i.e. the list1 contains nearest adjacent for the central point of the circle as in step 6). If a point p is covered by more than one circle then $p.distance$ keeps its distance to the nearest circle (step 8), also $p.circle$ keeps the identification of the nearest circle as in step 9. The function assign all points whose distances from this center is larger than Rad and less than $Rad * 1.5$ to the list2 as in step 12 (i.e. points which lie between the solid and dashed circle are the nearest far adjacent for the central point of the circle as shown in Fig. 3). For each point in list1, the algorithm keeps the distance to the center of the nearest circle and the identification number of that circle as in step 8 and step 9. The center of the next circle is the point whose distance is larger than $Rad * 1.5$ and less than or equal to $Rad * 2$ from the center of current circle to ensure from the existence of overlapping (this is shown in step 14). This process continues till all points are covered. Steps from 20 to 22 search for uncovered points remaining to cover them by creating new circles.

Thus every circle contains two list; the first list contains all points inside the solid edge (nearest adjacent for a central point as in Fig. 3), the other list contains all points outside the solid edge and inside the dashed edge (nearest far adjacent for a central point as in Fig. 3). Only, the algorithm uses the points inside the solid edge to find maximum of (minimum pairs wise distance) to compute the optimal value of Eps, that will be used by the DBSCAN algorithm in the next stage.

When the algorithm retrieves the neighbors of a point, it directly goes to the best (nearest) circle that covers the point (a point may be covered by more than one circle), computes how far is it from all the points in that circle, and returns points at distances less than or equal to Eps. If the distance of that point to the solid edge is less than Eps then the algorithm computes the distances between that point and all the points contained in dashed circle (points in list1 and list2), retrieves the points in Eps distance, and the point is assigned to current cluster if it is a core point, or assigned noise temporally.

V. THE PROPOSED ALGORITHM

The proposed algorithm merges ideas from many algorithms. It is based on DBSCAN while we try to solve the problems of Eps and R*-tree. In this section, the R*-tree is discussed, how the value of Eps is computed, and the computation of the radius of all data space.

A. R*-tree and overlapped circles

What is R*-tree?. What is the problem of it?. The R*-tree [1] generalizes the 1-dimensional B-tree to d-dimensional data spaces, specifically an R*-tree manages d-dimensional hyper rectangles instead of 1-dimensional numeric keys. An R*-tree may organize objects such as polygons using minimum bounding rectangles (MBR) as approximations as well as point objects as a special case of rectangles. The leaves store the MBR of data objects and a pointer to the exact geometry. Internal nodes store a sequence of pairs consisting of a rectangle and a pointer to a child node. These rectangles are the MBR's of all data or directory rectangles stored in the subtree having the referenced child node as its root (Fig. 5).

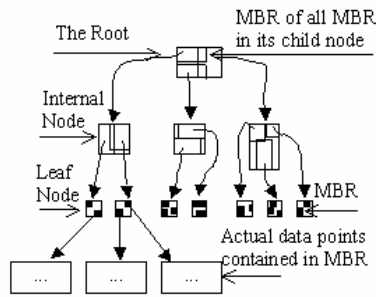


Fig. 5 Sample R*-tree

To answer a region query, we start from the root[4], and the set of rectangles intersecting the query region is determined and then their referenced child nodes are searched until the data pages are reached (i.e. more than one path is searched) and the search space become large. Since the overlap of the MBR's in the directory nodes grows with increasing dimension d , the R*-tree is efficient only for moderate values of d [4].

We partition the data space into overlapped circular (sphere or hyper sphere) regions such that the radius of each circle is larger than the expected Eps. We use these overlapped circles to answer a region query instead of R*-tree. This idea comes from Fig. 1, the overlapped circle of the same radius may be used to cover all data space with respect to the radius of circle is greater than the expected Eps. Some data points may be belonging to more than one circle but we use the nearest circle to retrieve it's neighborhood using Euclidean distance. This search time is better than that of R*-tree, since in R*-tree to answer a region query, we start from the root, and the set of rectangles intersecting the query region is determined and then their referenced child nodes are searched until the data pages are reached (i.e. more than one path is searched) and the search space become large. In the other hand, in the proposed algorithm say k of circles cover all data space. Thus, the search requires $O(mn)$, where n is the number of data points, $m=n/k$ is very small compared with n . For each point we keep the nearest circle center and the distance to it. The radius of the circle depends on the radius of all data space.

B. Computation of the Radius of the Data Space and Eps

How does the proposed algorithm determine the radius of all data space? The proposed algorithm accumulates all data points in a single cluster called cluster feature (CF); CF is a data structure summarizing information about all points in the dataset, $CF = (n, \bar{LS})$, where \bar{LS} is the linear sum of the n data points i.e. $\bar{LS} = \sum_{i=1}^n \bar{x}_i$, n is the number of points in dataset, \bar{x}_i is a d -dimensional data point. The center of the all data points in the dataset is $\bar{x}_0 = \bar{LS} / n = \sum_{i=1}^n \bar{x}_i / n$. The radius of the data space is given

by $R = \sqrt{\sum_{i=1}^n (\bar{x}_i - \bar{x}_0)^2 / n}$. So, R is the radius of the circle

(sphere) that contains all data points in our dataset. we compute the area of that circle from the relation "area = $3.14 * R^d$ " (we refer to this area as circular area), then we compute the area from other view, by using minimum bounding rectangle that contains all data points in the dataset,

"area = $\prod_{i=1}^d L_i$ ", where d is the dimension of the points, L_i is

the length of dimension i , which is equal to the difference between the maximum and the minimum value in dimension i (we refer to this area as rectangular area). Figure 6 shows the circular area and rectangular area of some data points in 2 dimension.

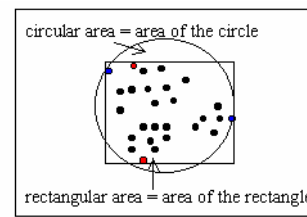


Fig. 6 Circular and Rectangular area of data space

In Fig. 6, the length of the first dimension is determined by the two blue points (x dimension). The length of the second dimension is determined by the two red points (y dimension).

The proposed algorithm partitions data space into overlapped circles. The radius of these circles should be depends on the dimension of data space, since as the dimension increase the data points will be more sparse. Also, the radius should be depends on the area of data space, but which area we can use? Experimentally the ratio between the two area is the best, so ratio area = (circular area / rectangular area) or the inverse i.e. ratio area = (rectangular area / circular area). The ratio area should be less than or equals to one i.e. $0 < \text{ratio area} < 1$. Experimentally the best relation for the radius of the overlapped circles is $Rad = d * \text{ratio area} + \text{ratio area} / 2$, where d is the dimension of the data space, ratio area is the ratio between circular area and rectangular area or the inverse. As the dimension increase the radius of the overlapped circles increase. Also, as the difference between the two areas decrease the ratio area increase and the radius of the overlapped circles increase. Always, the radius of the overlapped circles is greater than the expected Eps.

Here, we are ready to apply DBSCAN, but we will use circles that cover all data space to calculate the optimal value for Eps. To do this, for each point in the solid circle we find the distance to its nearest neighbor, then we keep the distance between the far nearest pairs, we perform this process for all circles, this process is equivalent to distances matrix update in the single link algorithm. The following example explain the idea.

Example 1: suppose one of the overlapped circles contains the following six points in 2 dimensional space.

TABLE I
SIX POINTS IN 2 DIMENSION

	Attribute1	Attribute2
P1	1	1.1
P2	1.25	0.9
P3	1	1.25
P4	1.23	1.22
P5	1.12	1.5
P6	1.13	0.7

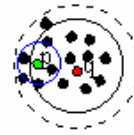


Fig. 7 Neighbors of the green point wrt. Eps

Then all pairs wise distances are in Table II (distances matrix) these distances are calculated by using the next equation.

$$d(p_i, p_j) = \sqrt{\sum_{k=1}^d (p_{i,k} - p_{j,k})^2}$$

Note that, all six points are in list1 i.e. points are inside the solid edge only.

TABLE II
PAIRS WISE DISTANCES (DISTANCES MATRIX)

	P1	P2	P3	P4	P5	P6
P1	0.000	0.320	0.150	0.259	0.418	0.421
P2	0.320	0.000	0.430	0.321	0.614	0.233
P3	0.150	0.430	0.000	0.232	0.277	0.565
P4	0.259	0.321	0.232	0.000	0.301	0.530
P5	0.418	0.614	0.277	0.301	0.000	0.800
P6	0.421	0.233	0.565	0.530	0.800	0.000

From Table II we find the nearest neighbor for each point is as follow:-

- P1 & P3 at 0.150
- P2 & P6 at 0.233
- P3 & P1 at 0.150
- P4 & P3 at 0.232
- P5 & P3 at **0.277**
- P6 & P2 at 0.233

The nearest pair are p1 and p3 at 0.150. The far nearest pair are p5 and p3 at **0.277**.

For this circle, the maximum distance between the nearest pair is 0.277 that is between p3 and p5. So the algorithm keeps this distance, then the algorithm take the next circle, this process is performed for all overlapped circles. Then Eps will be the average of these distances. i.e. $Eps = \sum max_i / k$, where k is the number of created overlapped circles and max_i is the distance between the far nearest pair of points in circle number i .

We use Eps as calculated before to overcome the presence of outliers. In the proposed algorithm the Minpts is fixed to 3 by the experiments.

Why Minpts is fixed to 3? 3 points is the best number for Minpts according to the method that we used to determine the value of Eps, if you examine the points in example 1, you find that all the six points are outliers

When do we use the points in list2 (points between solid and dashed circle)? To answer this question, look at the following Fig. 7.

Suppose that Rad = 1, Eps = 0.4, distance between the green point p and the red point q (the center) = 0.8, then the distance between the green point p to the solid edge = 0.2 and this distance is less than Eps. So, in this case we calculate the distance between the green point p and all points in dashed circle to find its neighbors wrt. Eps (points inside the blue circle (the smallest circle)).

After determination of Eps we apply the basic process of DBSCAN. We can summarize the basic process of the proposed algorithm in the following steps:

- 1- find the center of all data.
- 2- calculate the average radius of data (radius of circle that covers all data points).
- 3- divide the data into overlapped circular regions of the same radius such that this radius is larger than the expected Eps.
- 4- In each region, compute all pair wise distance (distance matrix as in Slink algorithm), find the maximum of (minimum distances). If we have k circles cover all data there will be k values for distance, we take the average value of these k distances to present the Eps of DBSCAN algorithm.
- 5- Apply DBSCAN algorithm on data using Eps obtained from step 4, to retrieve the neighbors of point p , only distances between point p and all points share the same region are calculated. This data structure is better than R*-tree, since in R*-tree more than one path is traversed and R*-tree works well with low dimensional data.

The proposed algorithm is the same as shown in Fig. 2, while Eps is not user input parameter, and overlapped circular region is used to answer a region query instead of R*-tree.

C. Complexity

As we discussed before, the proposed algorithm composed of two stage, in first stage, the algorithm creates k circular regions cover the data space, this requires $O(nk)$, where n is number of data point. To find the Eps the algorithm find pair wise distance in each region, in average each region contains m points, where $m = n/k$, thus this process takes $O(m^2k)$. So the time of first stage is $O(nk + m^2k)$. In second stage, the algorithm applies the DBSCAN, using the circular regions to answer query region. The search for the points in Eps distance from a random point requires $O(m)$, thus the DBSCAN requires $O(nm)$. Hence the total time complexity of the proposed algorithm is $O(nk + m^2k + nm)$, where k is the number of circles cover the data space, m is the average number of points in each circle, $m = n/k$ and n is number of points in the data set.

VI. EXPERIMENTAL RESULTS

We evaluated the proposed algorithm on several different real and synthetic datasets. We compared our results with that of DBSCAN algorithm in terms of the total quality of clusters, both algorithms produce the same result. Our experimental results are reported on PC 800 MHZ, 128 MB RAM, 256 KB cache. we give a brief description of the datasets used in our algorithm evaluation. Table III shows some characteristics of the datasets. The real data used in the experiments were taken from <http://www.cs.utoronto.ca/~delve/data/datasets.html> (abalone dataset), <http://lib.stat.cmu.edu/datasets/> (wind dataset) and earthquake dataset is send to me by the author of [4].

TABLE III
CHARACTERISTIC OF THE DATASETS

Datasets	Number of records	Number of attributes	Type of dataset
Earthquake	2049	2	real
Abalone	4177	7	real
Wind	6574	15	real
Db1	10000	2	synthetic

In Table IV, we present the results obtained from the proposed algorithm, present the optimal value for Eps that is always less than the radius of circle (Rad), number of circles present total number of circle that cover all data points. This number is equivalent to the number of leaf nodes in R*-tree. But in case of R*-tree we can not directly reach the required leaf node. We may also need to reach to more than one leaf node, this problem grows as the dimension of data grow. In the proposed algorithm only one circle is directly reached, this process saves time.

Also in Table IV, we present the results obtained from the DBSCAN algorithm compared with that of the proposed algorithm. The input values for Eps parameter is the same as in the proposed algorithm. Comparing the results of the proposed algorithm and DBSCAN algorithm, both algorithm produce the same results nearly.

TABLE IV
RESULTS OF THE PROPOSED AND DBSCAN ALGORITHM

Datasets	Radius of all data	Number of circles	Rad	Eps	Time (sec.)		Number of clusters		Number of noise points	
					proposed	DBSCAN	proposed	DBSCAN	proposed	DBSCAN
Earthquake	0.599	27	0.317	0.124	1	3	6	6	35	35
Abalone	0.489	30	0.245	0.152	14	24	2	2	28	28
Wind	19.17	920	9.584	8.54	31	116	12	11	704	703
Db1	5.527	95	1.224	0.249	4	58	100	100	4	4

VII. CONCLUSION

In this paper, we presented a density based clustering algorithm require no input parameters. The proposed algorithm handles large data set efficiently and discovers any arbitrary shaped clusters of any size. This algorithm is based on partitioning the data into overlapped circular or hyper spherical regions and uses the best region to retrieve the neighborhood of any data point. Our experimental results

demonstrated the efficiency of the proposed algorithm. In future work we will study hierarchical clustering algorithms. We will propose a hierarchical clustering algorithm handle huge dataset that will not be depends on sample as in CURE algorithm.

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