

Computer Verification in Cryptography

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Abstract—In this paper we explore the application of a formal proof system to verification problems in cryptography. Cryptographic properties concerning correctness or security of some cryptographic algorithms are of great interest. Beside some basic lemmata, we explore an implementation of a complex function that is used in cryptography. More precisely, we describe formal properties of this implementation that we computer prove. We describe formalized probability distributions (σ -algebras, probability spaces and conditional probabilities). These are given in the formal language of the formal proof system Isabelle/HOL. Moreover, we computer prove Bayes' Formula. Besides we describe an application of the presented formalized probability distributions to cryptography. Furthermore, this paper shows that computer proofs of complex cryptographic functions are possible by presenting an implementation of the Miller-Rabin primality test that admits formal verification. Our achievements are a step towards computer verification of cryptographic primitives. They describe a basis for computer verification in cryptography. Computer verification can be applied to further problems in cryptographic research, if the corresponding basic mathematical knowledge is available in a database.

Keywords—prime numbers, primality tests, (conditional) probability distributions, formal proof system, higher-order logic, formal verification, Bayes' Formula, Miller-Rabin primality test.

I. INTRODUCTION

MATHEMATICAL proofs are often complex and hard to verify by their readers. Consequently, the application of formal proof systems are a useful approach in the area of verification. Formal and computer verification augment the traditional concept of software engineering by providing techniques that guarantee trustiness as well as correctness of software systems in a mathematical way. There are many possible applications of formal and computer verification like automotive, medical technology, information technology security and cryptography.

In cryptographic research, the concept of provable security is of great interest to ensure, whether a given cryptographic primitive can be regarded as secure, or its behaviour is not as intended. During the last ten years, the relevance of security proofs increased rapidly. Great improvements in that direction were achieved, e.g. by proving the Optimal Asymmetric Encryption Padding (OAEP) for the case of RSA ([4]), which is widely used in practice. Finding a proof of an encryption scheme is in general done by considering the computational complexity of the regarded cryptographic primitive. Proof constructions in the complexity theoretic approach consider reductions from the problem of breaking a cryptographic primitive to the underlying (probably) hard problem. These proof constructions are typically made by human work (paper

proof). Consequently, a proof about security of a cryptographic primitive may contain some incorrectness. An example for this consequence (with respect to OAEP) was described in [8].

Our idea related to the above described problem is to construct proofs of cryptographic primitives like encryption schemes in the complexity theoretic approach with computer support. In this case, formal computational complexity and formal probability theory has to be studied. For this purpose, we use the proof assistant Isabelle/HOL, which is successfully applied in the *Verisoft project*¹.

We describe formalized probability distributions (σ -algebras, probability spaces and conditional probabilities). These are given in the formal language of the formal proof system Isabelle/HOL. Moreover, we computer prove (a formalized version of) Bayes' Formula. Besides we describe an application of the presented formalized probability distributions to cryptography.

The correctness of the implementation of cryptographic functions such as encryption and digital signature is crucial for the security of computer systems. But in general, such functions are very complex which makes formal proofs very hard. This paper shows that computer proofs of complex cryptographic functions are possible by presenting an implementation of the Miller-Rabin primality test that admits formal verification. This test is a key ingredient of RSA and DSA key generation. We describe its computer verification with Isabelle/HOL. More precisely, it is possible to formally prove

$$k \text{ prime}, x < k, \gcd(x, k) = 1 \implies \text{prim}(x, k) = 1,$$

and

$$\text{All primality conditions hold for } k, x \Leftrightarrow \text{prim}(x, k) = 1.$$

for an implementation *prim*. Besides, this implementation will be part of a fully verified cryptography client.

In the remaining part of this paper, we give a computer proven version of Bayes' Formula and we show how a formal proof system can be used to prove the correctness of an implementation of a primality test. As a result, a verified applicable function and formally proven properties, that extend the used database, are received. For this purpose, we give a new functional/logical description (and implementation) of the well known Miller-Rabin primality test. This description is written down and (interactively) verified in the formal proof language of Isabelle/HOL. Moreover, this implementation can be integrated into a cryptographic client that is implemented and verified in the used formal proof language.

Altogether we prove that formal verification with computer support of cryptographic relevant algorithms improves cryptographic practice by providing correct implementations, as well

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¹Verisoft is a German industrial research project, where formal mathematics, as well as computer science are applied to engineering.

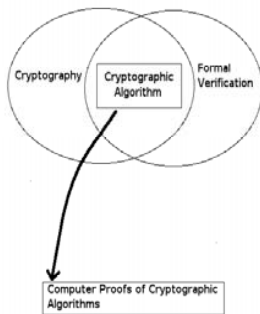


Fig. 1. Cryptography and Formal Verification

as cryptographic research by extending a formal database for the purpose of cryptography. Both aspects are illustrated below (Figure 1).

This paper is organized as follows: We start in II with a description of the used formal proof system. In III we explore Bayes' Formula with a formal proof system, and in IV we apply formalized probabilities in cryptography. Moreover, in V we give some mathematical definitions, that describe a basis of a formal description and verification of the Miller-Rabin algorithm. A formal verification of the Miller-Rabin algorithm is described in VI. Besides, this paper contains formal definitions of functions describing the Miller-Rabin algorithm, computer verified properties of these functions, and further explanations. In VII some conclusions, as well as some comments on future work are given.

II. THE ISABELLE SYSTEM

In the following, we give a short description of the verification tool Isabelle/HOL. For the reason of a better understanding of our formal version of (conditional) probability distributions, Bayes' Formula, cryptographic application of probability, and implemented version of the Miller-Rabin algorithm, which are formulated and verified in the formal language of Isabelle/HOL, we want to explain the main structures of the Isabelle proof system. It is a proof assistant for higher-order logic, which can be used for interactive proof constructions, formal specifications, as well as verification in higher-order logic and functional programming.

The formal language that consists of higher-order logic and functional programming, is used to give definitions and lemmata, which are based on a large database. These definitions and (proven) lemmata can be used to prove further lemmata and theorems, which results in an augmented database for the purpose of building up new theories (compare to Figure II).

More information about Isabelle/HOL (that is successfully applied in the *Verisoft project* (<http://www.verisoft.de>)) are given in [7], which describes constructions with this tool. A further useful reference is [9]. There, parts of the large database are mapped. Besides [9] contains other references about Isabelle/HOL.

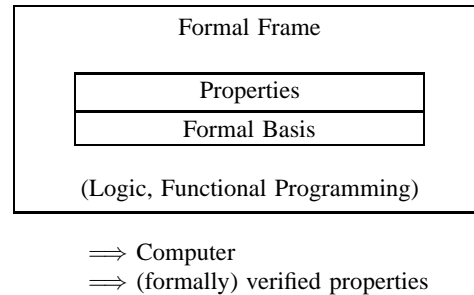


Fig. 2. Construction pattern in Isabelle/HOL described in II

III. BAYES' FORMULA

In the following, we give a computer proven version of Bayes' Formula that is formalized in the formal proof language of the used formal proof system. Bayes' Formula is a useful lemma to compute probabilities (for example computation of probabilities in cryptography). That means, a formalized version of this lemma is useful for formal proofs in cryptography.

A. Introduction

In order to computer prove Bayes' Formula we explore formal definitions of basic mathematical concepts (σ -algebra, probability space, conditional probability).

Review: A probability space (Ω, \mathcal{A}, P) is built up of a set Ω (set of results), a σ -algebra \mathcal{A} of Ω (system of possible events) and a probability distribution $P : \Omega \rightarrow [0, 1]$. A σ -algebra \mathcal{A} of Ω is a system of carriers from Ω ($\mathcal{A} \subseteq \mathcal{P}(\Omega)$). In the following these components of a probability space (i.e. their formal description) are presented (compare [1], [3], [5], [6] for more mathematical background).

1) *Formalized Mathematical Background:* We formalized a σ -algebra \mathcal{A} of a set E as a predicate `sigma_algebra` with a set of a fixed type and a set of sets of the same type as arguments. The declaration and definition of this predicate (constant) in the formal language of the Isabelle system is given as follows.

```
consts sigma_algebra :: "[ 'a set, 'a set set ] => bool";
```

defs

```
sigma_algebra_def:
  " sigma_algebra E A == E : A  $\wedge$  (ALL K : A. E-K : A)
   $\wedge$  (ALL B. ALL (i::nat). (B i) : A  $\longrightarrow$  ( $\bigcup$  i. B i) : A)";
```

This constant describes a predicate, which holds, if and only if the second argument A is a σ -algebra of the set E (first argument).

Moreover, we described a probability space through the declaration and definition of another predicate.

consts

```
pr_space :: "[ 'a set, 'a set set, ('a set => real) ] => bool";
```

defs

```
pr_space_def:
  "pr_space E F Pr == (sigma_algebra E F  $\wedge$  E  $\neq$  {})
```

$$\wedge (\text{ALL } A : F. 0 \leq (\text{Pr } A))$$

$$\wedge (\text{Pr } E) = 1 \wedge (\text{ALL } A : F. \text{ALL } B : F. ((A \cap B = \{\})$$

$$\longrightarrow (\text{Pr } (A \cup B)) = (\text{Pr } A) + (\text{Pr } B)))));$$

This predicate holds, if and only if its arguments form a probability space.

Besides we formalized the term of conditional probability in the formal language of Isabelle. Therefore the following predicate describes conditional probabilities.

consts cond_pr :: "[’a set, ’a set set, (’a set => real), (’a set, ’a set] => real] => bool";

defs

cond_pr_def:

”cond_pr E Alg Pr CoPr == pr_space E Alg Pr

$\wedge (\text{ALL } A : \text{Alg}. \text{ALL } Z : \text{Alg}. 0 < \text{Pr } Z$

$\longrightarrow \text{CoPr } A \ Z = \text{Pr } (A \cap Z) / \text{Pr } Z)$ ”;

2) *Bayes’ Formula*: Besides a formal definition of conditional probabilities, we have formalized some lemmata and Bayes’ Formula.

A σ -algebra of E , $A_0 \in \mathcal{A}$, $A_1 \in \mathcal{A}$, and $0 < \text{Pr}(A_0)$, $0 < \text{Pr}(A_1)$

$$\implies \text{Pr}(A_0|A_1) = \frac{\text{Pr}(A_1|A_0) \cdot \text{Pr}(A_0)}{\text{Pr}(A_1)}.$$

Our formalized version of Bayes’ Formula within the proof language of Isabelle is given below.

lemma: ”[[cond_pr E A Pr CoPr;

$A_0 \in \mathcal{A}; A_1 \in \mathcal{A}; 0 < \text{Pr } A_0; 0 < \text{Pr } A_1 \] \implies$

$(\text{CoPr } A_0 \ A_1) = (\text{CoPr } A_1 \ A_0) * (\text{Pr } A_0) / (\text{Pr } A_1)$ ”;

IV. CRYPTOGRAPHIC APPLICATION

In the lines below, we describe an application of the above presented formal framework to cryptography. Besides the explored application of formalized probability theory to cryptography, there are further possibilities of application (the term of *perfect secrecy* and related proofs provide an interesting area of application in provable security, but a deeper discussion won’t be given at this place).

A. Bellare-Rogaway

In [2] a well known proof technique is described (Bellare-Rogaway model, 1993). The following lines give a definition of semantical security against chosen ciphertext attack (IND-CCA).

Definition 4.1: IND-CCA

A given encryption scheme with security parameter z is IND-CCA in the Bellare-Rogaway model, if the following holds for all polynomial time algorithms A and for all $f \in \mathbf{Z}$ with $f > 0$.

For f there is x , for all $z > x$:

IND-CCA: $|2 \cdot \text{Pr}(A \text{ finds correct bit}) - 1| = \text{Adv}(A) < \frac{1}{z^f}$,

where the adversary is represented by A .

In this definition, the term $\text{Pr}(A \text{ finds correct bit})$ specifies the probability an algorithm A has to learn, whether a certain

bit is equal to 0 or 1, if it attacks a cryptosystem with security parameter z .

Below we want to define the term of IND-CCA in the formal language of Isabelle/HOL.

B. Formal Definition of IND-CCA

The following formal definition of IND-CCA depends on our version of formalized probability theory (which is based on the theory *Main* and on theories about real numbers and sets).

consts IND_CCA ::

”[’a set, ’a set set, (’a set => real), ’a set] => bool”;

defs IND_CCA_def:

”IND_CCA E Alg Pr A == pr_space E Alg Pr

$\wedge \forall (f::\text{int}). 0 < f \longrightarrow (\exists (x::\text{int}). 0 < x \longrightarrow$

$(\forall (z::\text{int}). x < z \longrightarrow |2 * (\text{Pr } A) - 1| < 1/z^f))$ ”;

In the lines above IND_CCA represents a predicate, that describes the dependencies between the probability of A (represented by the term $\text{Pr } A$) and the value of z^f .

V. MILLER-RABIN ALGORITHM

In this section we present the Miller-Rabin primality test. That presentation will be the basis of the Isabelle/HOL implementation described in the next section.

A. Primality or Compositeness

In general, a primality test is an algorithm that reveals, whether a given integer number k is prime or composite. In this context, a concrete prime factor decomposition of a composite number is not discussed.

Public-key cryptography often uses large prime numbers, consequently efficient algorithms are needed. A well known algorithm is division by possible factors, but because of its complexity ($O(\sqrt{k})$) this algorithm is not used for large numbers. More efficient algorithms are the *Fermat algorithm* and the *Miller-Rabin algorithm* ([3] or [10]), but their results are incorrect with a certain (small) probability.

B. Miller-Rabin Algorithm – Overview

In order to test an odd positive integer k for primality, the Miller-Rabin test proceeds as follows. As a precomputation the Miller-Rabin test determines the decomposition

$$k - 1 = 2^z \cdot v,$$

where $z \in \mathbf{N}_{\geq 1}$ and v is an odd positive integer. Functions that compute these two numbers are given in VI.

The Miller-Rabin test looks for witnesses that attest the compositeness of k . Such a witness is a number $x \in \{0, \dots, k-1\}$, $\text{gcd}(x, k) = 1$, such that the following property holds.

Theorem (1).

For $k \in \mathbf{N}$ prime, $z = \max\{r \in \mathbf{N} : 2^r \mid k - 1\}$, $v = (k - 1)/2^z$ ($k - 1 = 2^z \cdot v$), and $x \in \{0, \dots, k - 1\}$ randomly chosen with $\text{gcd}(x, k) = 1$ the following holds:

$$(1.1) \quad x^v \equiv 1 \pmod{k}$$

or there is a number $r \in \{0, \dots, z-1\}$ with

$$(1.2) \quad x^{2^r \cdot v} \equiv -1 \pmod{k}$$

If (1.1) and (1.2) do not hold for a given k , x is a witness for the compositeness of k . Theorem (1) can be used to implement an algorithm that looks for witnesses that attest the compositeness of k (compare [3]).

C. Miller-Rabin Algorithm – Implementation

In this section we describe an implementation of the Miller-Rabin test. We use Theorem (1) to implement Algorithm (1) that looks for witnesses that attest the compositeness of k . If no such witness is found, primality of k is assumed.

Algorithm (1) uses the precomputation introduced above, i.e. input of Algorithm (1) are an odd positive integer k (the primality/compositeness of k is tested), a randomly chosen number $x \in \{0, \dots, k-1\}$ with $\gcd(x, k) = 1$, and the decomposition $k-1 = 2^z \cdot v$. Output of Algorithm (1) is *composite* or *prime*. We use the two functions $\text{prim}_1 : \mathbf{N}^4 \rightarrow \mathbf{N}$ and $f_1 : \mathbf{N}^3 \rightarrow \mathbf{N}$ that are based on Theorem (1) for a search for witnesses. These functions are defined as follows.

$$\text{prim}_1(x, k, z, v) = \begin{cases} \text{if } x^v \equiv 1 \pmod{k} \text{ then } 1 \\ \text{else } f_1(z-1, x^v \pmod{k}, k) \end{cases}$$

uses f_1 , where

$$f_1(z, x, k) = \begin{cases} \text{if } 0 < z \text{ then} \\ \quad \text{if } x \equiv -1 \pmod{k} \text{ then } 1 \\ \quad \text{else } f_1(z-1, x^2, k) \\ \text{else } 0. \end{cases}$$

prim_1 and f_1 are applied to x , k , z , and v as described in the following lines.

Algorithm (1).

Input: $k \in \mathbf{N}$, $x \in \{0, \dots, k-1\}$ with $\gcd(x, k) = 1$, $k-1 = 2^z \cdot v$ with properties given above

Output: *composite* or *prime*

If $\text{prim}_1(x, k, z, v) = 1$ then return *prime*
else return *composite*

For some technical reasons we use the functions $\text{prim} : \mathbf{N}^4 \rightarrow \mathbf{N}$, $\text{prim}_0 : \mathbf{N}^4 \rightarrow \mathbf{N}$, $f : \mathbf{N}^3 \rightarrow \mathbf{N}$, and $\text{exp}_2 : \mathbf{N}^3 \rightarrow \mathbf{N}$ to implement and verify the Miller-Rabin test with Isabelle/HOL. These functions are defined as follows.

$$\text{prim}(x, k, z, v) = \begin{cases} \text{if } z = 1 \text{ then } \text{exp}_2(x, k, v) \\ \text{else } \text{prim}_0(x, k, z, v) \end{cases}$$

prim provides a case split, i.e. for $z = 1$ the function exp_2 computes a corresponding output value, while the function prim_0 looks for witnesses otherwise. This case distinction results from an abbreviation used in f that corresponds to f_1 .

$$\text{prim}_0(x, k, z, v) = \begin{cases} \text{if } x^v \pmod{k} \equiv 1 \text{ or } x^v \pmod{k} \equiv -1 \\ \text{then } 1 \\ \text{else } f(z-1, x^v \pmod{k}, k) \end{cases}$$

$$f(z, x, k) = \begin{cases} \text{if } 0 < z \text{ then} \\ \quad \text{if } x^v \pmod{k} \equiv -1 \text{ then } 1 \\ \quad \text{else if } x^v \pmod{k} \equiv 1 \text{ then } 0 \\ \quad \text{else } f(z-1, x^2 \pmod{k}, k) \\ \text{else } 0 \end{cases}$$

$$\text{exp}_2(x, k, v) = \begin{cases} \text{if } x^v \pmod{k} \equiv 1 \text{ or } x^v \pmod{k} \equiv -1 \\ \text{then } 1 \\ \text{else if } x^{2 \cdot v} \pmod{k} \equiv -1 \\ \text{then } 1 \\ \text{else } 0 \end{cases}$$

A version of the Miller-Rabin test that uses similar case distinctions is given in [10].

D. Miller-Rabin Algorithm – Correctness

Up to now we gave no justification why the described search for witnesses for compositeness of odd positive integer k should be successful. But there are sufficient many witnesses, if k is composite, what is explained in Theorem (2), that describes the probability of correctness.

Theorem (2).

For $k \in \mathbf{N}$, $k \geq 3$ odd, composite number, the set $\{1, \dots, k-1\}$ contains at most $(k-1)/4$ numbers x , where $\gcd(x, k) = 1$ and (1.1), (1.2) hold.

That means, if an integer number is composite, there are sufficient many witnesses to attest the compositeness of this number (compare [3]).

VI. VERIFICATION OF THE MILLER-RABIN ALGORITHM

Algorithm (1) (with prim instead of prim_1) is the main foundation of a formal description of the Miller-Rabin algorithm with Isabelle/HOL. This computer representation, that is based on functional programming and logic (of higher order), is a direct implementation of the functions prim , prim_0 , exp_2 and f given in (3).

A. Formalized Mathematical Properties

In the following, we give a formalized version of Theorem (1) that can be verified with Isabelle/HOL.

lemma "[(k:int) ∈ zprime; zgcd((x:int),k) = 1; 2 < k; x < k] ==> (∃ (z:num). z < (f_02 (number k)) ∧ (x^((2^z) * (mult2_value (number k)))) mod k = k - (1::int) | (x^(mult2_value (number k))) mod k = 1)";

In the lines below some Isabelle/HOL code examples are explained.

Remark: (Isabelle/HOL code)

- int: integer number datatype
- num: natural number datatype
- ^z: exponent z
- zprime: set of (integer) prime numbers
- zgcd: (integer) greatest common divisor
- f_02 and mult2_value compute a decomposition of $x-1 = 2^{(f_02x)} \cdot (\text{mult2_value}x)$, if $x \in \mathbf{N}$ is odd ($(f_02x) =$

$$\max\{r \in \mathbf{N} : 2^r \mid x - 1\}, (\text{mult2_value } x) = (x - 1)/2^{(\text{f_02 } x)}.$$

- function *number* describes the type change of an integer number z to a natural number corresponding to z

The presentation of this lemma not only shows that the implementation of the Miller-Rabin test given below was computer proven. Furthermore, the underlying mathematical theory can be verified with Isabelle/HOL.

B. Implementation

In this formal description, the functions *prim*, *prim₀*, *exp₂* and *f* are implemented by *recovered_value01* and *primality*, *recovered_value0*, *recovered_value2*, and *rep_operator*, respectively. Moreover, some auxiliary functions are used. These auxiliary functions are given below.

Auxiliary Functions:

We implemented the two auxiliary functions *f₀₂* and *mult2_value* as follows.

consts *f_02* :: "num \Rightarrow num";

defs

f_02_def: "f_02 x \equiv (GREATEST (z::num). (2^z z dvd (x-(1::num))))";

consts *mult2_value* :: "num \Rightarrow num";

defs

mult2_value_def: "mult2_value x \equiv (x-(1::num)) div (2^{f_02 x})";

These two functions compute a decomposition of $x - 1 = 2^{(\text{f_02 } x)} \cdot (\text{mult2_value } x)$, if $x \in \mathbf{N}$ is odd. That means $(\text{f_02 } x) = \max\{r \in \mathbf{N} : 2^r \mid x - 1\}$ and $(\text{mult2_value } x) = (x - 1)/2^{(\text{f_02 } x)}$.

Functions describing the Miller-Rabin algorithm:

We directly implemented the functions *prim* (recovered_value01 and primality), *prim₀* (recovered_value0), *exp₂* (recovered_value2), *f* (rep_operator).

consts *rep_operator* :: "(int \times int \times int) \Rightarrow int";

recdef *rep_operator* "measure $(\lambda (z,x,k). (\text{number } z))$ "

rep_operator_def: "rep_operator (z,x,k) = (if (0 < z) then (if ((x² mod k) = (k - (1 :: int))) then (1 :: int) else (if ((x² mod k) = 1) then (0 :: int) else (rep_operator ((z - (1 :: int)), (x² mod k), k)))) else 0)"

(hints recdef_simp: number_lemma);

consts *primality* :: "[int, int, int, int] \Rightarrow int";

defs

primality_def: "primality x k z v \equiv (if ((x^{number v} mod k) = 1) | ((x^{number v} mod k) = (k - (1 :: int))) then (1 :: int) else (rep_operator ((z - (1 :: int)), (x^{number v} mod k), k))))";

consts *recovered_value0* :: "int \Rightarrow int \Rightarrow int";

defs

recovered_value0_def: "recovered_value0 x k \equiv primality x k (int (f_02 (number k))) (int (mult2_value (number k)))";

consts *recovered_value2* :: "int \Rightarrow int \Rightarrow int";

defs

recovered_value2_def: "recovered_value2 x k \equiv (if ((x^{number v} mod k) = 1) | ((x^{number v} mod k) = (k - (1 :: int))) then (1 :: int) else if (x^{number v} mod k) = k - (1 :: int) then (1 :: int) else (0 :: int))";

consts *recovered_value01* :: "int \Rightarrow int \Rightarrow int";

defs

recovered_value01_def: "recovered_value01 x k \equiv (if (f_02 (number k) = 1) then (recovered_value2 x k) else (recovered_value0 x k))";

In this description we used the function *number*. This function describes the type change of an integer number z to a natural number corresponding to z .

C. Verification

The new functional and logically compiled description as well as its implementation (formal and machinery description) of the Miller-Rabin algorithm given above, illustrate that formalized mathematical knowledge can improve verification of correctness of primality tests or other (cryptographic) algorithms. This improvement is reached by using Isabelle/HOL as computer verification tool. We interactively proved the following main results that we give in a mathematical description).

Theorem (3). $x, k \in \mathbf{Z}$, $k \in \text{Prime}$, $\text{gcd}(x, k) = 1$, $2 < k$, $x < k$, $0 < \text{f_02}(k)$, $0 < \text{mult2_value}(k) \implies \text{prim}(x, k) = 1$.

(If an integer number k is prime and some preconditions hold, the function *prim* outputs 1 (what indicates that the Miller-Rabin test is passed).)

Theorem (4). $x, k \in \mathbf{Z}$, $x < k$, $2 < k$, $0 < \text{f_02}(k)$, $0 < \text{mult2_value}(k) \implies ((\exists z \in \mathbf{N}. z \leq \text{f_02}(k) \wedge x^{2^z} \cdot \text{mult2_value}(k) \bmod k = k - 1) \vee (x^{\text{mult2_value}(k)} \bmod 1) = (\text{prim}(x, k) = 1))$.

(Under some preconditions, $((\exists z \in \mathbf{N}. z \leq \text{f_02}(k) \wedge x^{2^z} \cdot \text{mult2_value}(k) \bmod k = k - 1) \vee (x^{\text{mult2_value}(k)} \bmod k = 1))$ is equivalent to the fact that *prim* outputs 1.)

In the following, we illustrate a computer verification of these results by giving some examples of properties we proved with Isabelle/HOL.

After the implementation of the above introduced functions describing the Miller-Rabin algorithm, we started the proof construction with interactive proofs of some easy lemmata.

Example: (For a better reading we skipped the data types (e.g.int) of the variables)

lemma "[| 2 < k; k-1 = 2^{(number z) * v}; x < k; x^{(number v) mod k} = k-1 |] \implies (primality x k z v) = 1";

apply (unfold primality_def);
 apply (auto);
 done;

(lemma with computer proof script)

Many lemmata of that kind and further lemmata of increasing difficulty lead us to a proof of the following property.

lemma "[| 0 < z; $\exists i. (\forall j. j < i \longrightarrow (x^{(2^j)} \bmod k)^2 \bmod k \neq k-1) \wedge (\forall j. j < i \longrightarrow (x^{(2^j)} \bmod k)^2 \bmod k \neq 1) \wedge ((x^{(2^i)} \bmod k)^2 \bmod k = k-1) \wedge 0 < i \wedge 0 < z - (\text{int } i) \wedge \forall j. j < i \longrightarrow 0 < z - (\text{int } j) \rceil \rceil \implies (\text{rep_operator } (z, (x \bmod k), k)) = 1$ ";

Consequently, we were able to prove the following lemma.

lemma "[| 2 < k; k-1 = 2^(number z) * v; x < k; 0 < z-1; (x^(number v) mod k) = k-1) \vee (x^(number v) mod k) = 1) \vee ($\exists i. (\forall j. j < i \longrightarrow (x^{(2^j)} \bmod k)^2 \bmod k \neq k-1) \wedge (\forall j. j < i \longrightarrow (x^{(2^j)} \bmod k)^2 \bmod k \neq 1) \wedge ((x^{(2^i)} \bmod k)^2 \bmod k = k-1) \wedge 0 < i \wedge (\text{int } i) < z-1) \rceil \rceil \implies (\text{primality } x \ k \ z \ v) = 1$ ";

For a complete formal verification, we introduced the following facts to the formal proof system:

- "(k::int) - (1::int) = ((2^(number (int (f_02 (number k)))) * (int (mult2_value (number k))))"
- "[| 2 < k; x < k; 0 < f_02 (number k); 0 < mult2_value (number k); recovered_value01 x k = 1 |]"

Fact 1 was used to proof the following properties.

lemma "[| k \in zprime; zgcd(x,k) = 1; 2 < k; x < k; 0 < f_02 (number k)-1; 0 < mult2_value (number k) |]" \implies recovered_value0 x k = 1";

(correctness of auxiliary function recovered_value0)

lemma "[| k \in zprime; zgcd(x,k) = 1; 2 < k; x < k; f_02 (number k) = 1 |]" \implies recovered_value2 x k = 1";

(correctness of auxiliary function recovered_value2)

lemma "[| k \in zprime; zgcd(x,k) = 1; 2 < k; x < k; 0 < f_02 (number k); 0 < mult2_value (number k) |]" \implies recovered_value01 x k = 1";

(correctness of recovered_value01 (implementation of prim) (Theorem (3))

More (proven) lemmata and Fact 2 were used to verify the property below.

lemma "[| x < k; 2 < k; 0 < f_02 (number k); 0 < mult2_value (number k) |]" $\implies ((\exists z. z < (f_02 (number k)) \wedge (x^{(2^z)} * (\text{mult2_value } (\text{number } k)))) \bmod k = k-1) \vee (x^{(\text{mult2_value } (\text{number } k))} \bmod k = 1) = (\text{recovered_value01 } x \ k = 1)$ ";

(Theorem (4))

Altogether we implemented a functional/logic version of the Miller-Rabin primality test, that means we are able to computer verify the correctness of our algorithm. Besides, the used functions can be applied to integer numbers, what results in a applicable computer program for primality tests. In the following lines this result is illustrated.

Example: Application of ML code of

recovered_value0 x k

yields 0 or 1 corresponding to the values of $x, k \in \mathbf{Z}, x < k$.

VII. CONCLUSION

We explored formalized probability distributions (σ -algebras, probability spaces as well as conditional probabilities). These are given in the formal language of the formal proof system Isabelle/HOL. Moreover, we computer proved a formalized version of Bayes' Formula. Besides, we described an application of the presented formalized probability distributions to cryptography. Furthermore, this paper showed that computer proofs of complex cryptographic functions are possible by presenting an implementation of the Miller-Rabin primality test that admits formal verification. In this paper we gave a new description of a well known algorithm, that means we implemented the Miller-Rabin algorithm with computer support. Our new functional/logic implementation is computer verified and applicable as ML code. Moreover, we are able to computer verify the correctness of our algorithm. This paper is a further step towards a formalized basis of cryptographic algorithms. In cryptographic applications large prime numbers are often needed and therefore efficient correct algorithms for primality tests must be implemented and verified. Formal verification is a appropriate account to achieve correct algorithms. Moreover, cryptographic research can gain profit from formalized knowledge and formal verification. Formal proofs of security may support verification of cryptographic protocols or algorithms like encryption schemes. For that purpose, further mathematical knowledge must be implemented in a formal proof system.

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